## EyeDB Object Query Language

Version 2.8.8

Copyright © 1994-2008 SYSRA

Published by SYSRA
30, avenue Général Leclerc
91330 Yerres - France
home page: http://www.eyedb.org

## Contents

1 Introduction ..... 5
2 Principles ..... 5
3 OQL vs. ODMG 3 OQL ..... 5
4 Language Concepts ..... 7
5 Language Syntax ..... 8
5.1 Terminal Atom Syntax ..... 8
5.2 Non Terminal Atom Production ..... 10
5.3 Keywords ..... 11
5.4 Comments ..... 11
5.5 Statements ..... 12
5.6 Expression Statements ..... 12
5.7 Atomic Literal Expressions ..... 14
5.8 Arithmetic Expressions ..... 14
5.9 Assignment Expressions ..... 18
5.10 Auto Increment \& Decrement Expressions ..... 19
5.11 Comparison Expressions ..... 20
5.12 Logical Expressions ..... 24
5.13 Conditional Expression ..... 24
5.14 Expression Sequences ..... 25
5.15 Array Deferencing ..... 25
5.16 Identifier Expressions ..... 28
5.17 Path Expressions ..... 32
5.18 Function Call ..... 33
5.19 Method Invocation ..... 34
5.20 Eval/Unval Operators ..... 37
5.21 Set Expressions ..... 38
5.22 Object Creation ..... 40
5.23 Object Deletion ..... 42
5.24 Collection Expressions ..... 42
5.25 Exception Expressions ..... 49
5.26 Function Definition Expressions ..... 49
5.27 Conversion Expressions ..... 51
5.28 Type Information Expressions ..... 54
5.29 Query Expressions ..... 55
5.30 Miscellenaous Expressions ..... 59
5.31 Selection Statements ..... 60
5.32 Iteration Statements ..... 61
5.33 Jump Statements ..... 63
5.34 Function Definition Statements ..... 64
6 Quick Reference Manual ..... 68
6.1 Builtin and Library Functions and Methods ..... 68
6.2 Special Variables ..... 71
6.3 The eyedboql Tool ..... 71
6.4 The Standard Library Package ..... 75
6.5 OQL Quick Reference Card ..... 89

## The Object Query Language

## 1 Introduction

In this chapter, we present the EyeDB Object Query Language which supports the EyeDB object model. It is based on the ODMG 3 Object Query Language OQL

We first describe the design principles of the language in Section 2, then we present in Section 3 the main differences between EyeDB OQL and ODMG OQL. The language concepts are presented in Section 4. In Section 5, we introduced the language syntax. A quick reference manual of OQL is given in Section 6.

In this chapter, OQL denotes the EyEDB Object Query Language while the standard ODMG 3 Object Query Language will be denoted as ODMG OQL.

## 2 Principles

The principles of OQL are close to the principles of ODMG OQL introduced in the book Object Database Standard, ODMG 3 by Rick Cattell and al.

Our design is based on the following principles and assumptions:

- OQL relies on the EyeDB Object Model.
- OQL is based on ODMG OQL close to SQL 92. Extensions to SQL 92 concern object-oriented notions, like complex objects, object identify, path expressions, polymorphim, operation invocation and late binding. Extensions to ODMG OQL concern function definitions, selection statements, iteration statement, assignment operator (see Section 3).
- OQL provides high level primitives to deal with sets of objects, structures, lists and arrays.
- OQL is a functional language where operators can freely be composed as long as the operands respect the type system.
- OQL is computationally complete.
- OQL can be invoked from within programming languages for which an EyeDB binding is defined (currently C++ and Java). Conversely, OQL can invoke operations programmed in this language.


## 3 OQL vs. ODMG 3 OQL

OQL implements all the ODMG OQL functionalities with a few exceptions concerning the select clause. In order to accept the whole DML (Data Manipulation Language) query part of SQL as a valid syntax for ODMG OQL, ad-hoc constructions have been added to ODMG OQL each time SQL introduces a syntax that cannot be considered in the category of true operators.

For instance, the following construct is a valid ODMG OQL construct:
select p.name, salary from Professors p
This construct is not currently a valid OQL construct. The alternate valid OQL form is:

```
select struct(name: p.name, salary: p.salary) from Professors p
```

In the same way, ODMG OQL accepts SQL forms of the agregate operators min, max, count, sum and avg, for instance:

```
select count(*) from Persons
select max(e.salary) from Employees e
```

These constructs are not currently valid OQL constructs. The alternate valid OQL forms are:

```
count(select p from Persons p)
max(select e.salary from Employees e)
```

In the same way, the select $*$ clause is not currently implemented in OQL, neither the implicit select clause (i.e. without explicit identifier). For instance, the following constructs are not OQL valid constructs, although there are valid ODMG OQL constructs:

```
select * from Person
select name from Person
```

There is no alternate OQL valid form for the first construct. The alternate OQL valid forms for the second construct is:

```
select Person.name
select p.name from Person p
select p.name from Person as p
select p.name from p in Person
```

On the other hand, OQL provides a few extensions to ODMG OQL:

- assignment operators,
- four regular expression operators,
- selection statements, if/else,
- iteration statements, while, do/while, two forms of for,
- function definition statements, function,
- the eval and unval operators,
- identifier operators, isset, unset, refof, valof, \&, *, push, pop,
- exception management operators, throw,
- type information operators, classof, typeof
- miscellaneous operators, structof, bodyof
- builtin and library functions.

For instance, the following constructs are valid OQL constructs:

```
for (x in l)
{
    if (classof x != "Person")
        throw "Person type expected";
    if (x->name ~ "^john")
    {
            ok := true;
            break;
    }
}
function fib(n) {
    if (n < 2)
        return n;
    return fib(n-1) + fib(n-2);
}
for (n := 0, v := 0; n < 15; n++)
    v += fib(n);
function swap(x, y) {
    v := *x;
    *x := *y;
    *y := v;
}
```

```
i := "ii"; j := "jj";
swap(&i, &j);
function get_from(classname, attrname) {
    return eval "select x." + attrname + " from " + classname + " x";
}
names := get_from("Person", "name");
```

These extensions make OQL computationally complete.
Some of the ODMG OQL functionnalities or specificities are not yet implemented:

1. the group by/having operator,
2. the order by operator is more restrictive than in the ODMG specifications,
3. contrary to ODMG OQL, it is necessary to put parenthesis to call a function or method with takes no arguments,
4. contrary to ODMG OQL, the \| operator does means string concatenation. It is the logical or operator. This will be changed in a future version.

## 4 Language Concepts

The basic entity in OQL is the atom. An atom is the result of the evaluation of any expression. Atoms are manipulated through expressions.

Although OQL is not fully an expression language as some valid constructs, such as flow controls, are not expressions, the expression is the basic concept of OQL. The non-expression constructs, such as selection and iteration statements, control the flow (or evaluation) of expressions. An expression is built from typed operands composed recursively by operators.

OQL is a typed language: each atom has a type. This type can be an OQL builtin type or can be derived from the schema type declarations.

OQL binds the EyEDB object model by providing a mapping between OQL builtin types and the EyEDB object model types. As in the EYEDB object model, OQL supports both object entities (with a unique object identifier) and literal entities (with no identifier). The concept of object/literal is orthogonal to the concept of type, this means that any type may have object instances and literal instances. For instance, one can have literal or object integers, literal or object collections. For instance, the following constructs produces a literal integer:

```
1; // OQL interpreter generates a literal integer
first(select x.age from Person x); // database produces a literal integer
        // bound to an atom of type integer
```

while the followings produce respectively a Person object and a literal collection of Person objects:

```
first(select x from Person x);
select x from Person x;
```

An EyEDB object is always bound in OQL to an atom of type oid or of type object. A literal is bound to the corresponding direct type in OQL using a trivial mapping. For instance, a literal entity of the EyEDB type integer is bound to an OQL atom of the OQL type integer; while an object entity of the EyeDB type Person is bound to an OQL atom of type oid or object.

We introduce now the OQL builtin types and the way that they are generated. OQL includes 15 builtin types as follows:

> - integer

- string
- float
- char
- boolean
- identifier
- set
- bag
- array
- list

$$
\begin{aligned}
& \text { - struct } \\
& \text { - oid } \\
& \text { - object } \\
& \text { - null } \\
& \text { - nil }
\end{aligned}
$$

Some of these atoms can be expressed as terminals of the OQL grammar - for instance integer, float, string - others are generated using syntaxic constructions such as function calls or specific constructions - for instance list, bag, struct.

We will introduced first the syntax of the atoms which can be expressed as terminals, then the way to produce non terminal atoms.

## 5 Language Syntax

### 5.1 Terminal Atom Syntax

To express the syntax of terminal atoms, we use the standard regular expression notation.

## Integer Atom

Integers are coded on 64 bits.
The syntax for the integer type is one of the followings:

| $[0-9]+$ | decimal base |
| :--- | :--- |
| $0 \mathrm{x}[0-9 \mathrm{a}-\mathrm{fA}-\mathrm{F}]+$ | hexadecimal |
| $0[0-7]+$ | octal |

The domain for the integer type is as follows:

| Minimal Value | Maximal Value |
| :--- | :--- |
| -9223372036854775808 | 9223372036854775807 |

A few examples:

| 13940 | // integer expressed in the decimal base |
| :--- | :--- |
| $0 x 273 f 1$ | // integer expressed in the hexadecimal base |
| 0x273F1 | // integer expressed in the hexadecimal base |
| 0100 | // integer expressed in the octal base |

## Float Atom

The syntax for floating point atoms is one of the following regular expressions:

```
[0-9]+\.[0-9]+?
[0-9]+?\. [0-9]+
[0-9]+\.[0-9]+?(e|E)[+-]?[0-9]+([fF]|[1L])?
[0-9]+?\.[0-9]+(e|E)[+-]?[0-9]+([fF]|[lL])?
```

The domain for the float type is as follows:

| Minimal Value | Maximal Value |
| :--- | :--- |
| $4.94065645841246544 \mathrm{e}-324$ | $1.79769313486231570 \mathrm{e}+308$ |

A few examples:
1.
1.23
. 3
0.3039
$1 \mathrm{e}+10$
2.e+112
$1.2 \mathrm{e}-100$
. 234e-200
. 234e-200f
. $234 \mathrm{e}-200 \mathrm{~F}$

## String Atom

The syntax for the string type is as follows: \" ([~"] <br><br>")*""

The following escape sequences are interpreted:

| Escape Sequence | Name | ASCII Name |
| :---: | :---: | :---: |
| $\backslash \mathrm{a}$ | alert | BEL |
| $\backslash \mathrm{b}$ | backspace | BS |
| $\backslash \mathrm{f}$ | form feed | FF |
| $\backslash \mathrm{n}$ | newline | NL (LF) |
| $\backslash \mathrm{r}$ | carriage return | CR |
| $\backslash \mathrm{t}$ | horizontal tab | HT |
| \v | vertical tab | VT |
| \1 | backslash | $\backslash$ |
| \" | double quote | " |
| \' | single quote | , |
| $\backslash 000$ | octal number | \000 |

A few examples:
"hello"
"hello \"world\""
"this is a multi-lines $\backslash n t e x t \backslash n "$
"this text contains escape sequences: \007\v\f\n'"

## Char Atom

The syntax for the char type is one of the followings:
'ascii character'
,$\backslash[0-7+]$,
, $\backslash(\mathrm{x} \mid \mathrm{X})[0-9 \mathrm{a}-\mathrm{fA}-\mathrm{F}+]$,
, $\backslash a^{\prime}$
${ }^{\prime} \backslash \mathrm{b}$,
, \f,
${ }^{\prime} \backslash \mathrm{n}$,
${ }^{\prime} \backslash r$,
' $\backslash \mathrm{t}$ '
${ }^{\prime} \backslash \mathrm{v}$,

A few examples:
'a'
'b'
' $\backslash \mathrm{n}$,
'\a'
'\007'
' x 50 '
' x 5 F '

## Boolean Atom

The syntax for a boolean atom is one of the followings:
true
false

## Identifier Atom

The syntax for an identifier atom is as follows:
[a-zA-Z<br>\$_\#] [a-zA-Z<br>\$_0-9\#]*
This means that an identifier must start whith a letter, a " ", a " $\$$ " or a "\#" which may be followed by letters, digits, "_", "\$" and "\#" characters.

For instance, the following words are some valid identifiers:
a
alpha
beta1
alpha_beta
\$a
oql\$maxint
oql\#2
\$
\#
_1
Note that identifiers beginning by oql\$ or oql\# are reserved for special used by the interpreter.

## Oid Atom

The syntax for an oid is as following: [0-9]+: [0-9]+: [0-9]+: oid

Note that oid atoms are not typed directly by the user, but are produced by the database via the OQL interpreter. The following words are some syntaxically valid atom oids:
123.2.33373:oid
82727272.1.292828282: oid

## Object Atom

The syntax for an atom object is as following:
[0-9a-fA-F]+:obj

Note that object atoms are not typed directly by the user, but are produced by OQL interpreter. The following words are some syntaxically valid atom objects:

38383:obj
ea954:obj

## Null Atom

The null atom denotes an unitialized value. Its type depends on the context. It can denote a unitialized integer, float, char, string or oid.

The syntax for a null atom is one of the followings:
null
NULL

## Nil Atom

The nil atom denotes the empty atom. The syntax for a nil atom is as follows:
nil

### 5.2 Non Terminal Atom Production

The other atoms - sets, bags, arrays, lists and structs - are non terminal atoms. This means that they cannot be generated using a simple lexical construct.

## List, Set, Bag and Array Atoms

To construct a collection atom - list, set, bag or array -, one may use the function collection() where collection denotes the collection type, for instance:
$\operatorname{set}(1,2,3)$
list(1, "hello", "world")
$\operatorname{array}(2,3$, list (3893, -2 , 'a'), 22)
$\operatorname{bag}(2,2,3,4,5,12)$
This is the simple way to construct such atoms, but as any other atoms, a collection atom may be produced by the OQL interpreter as the evaluation of a complex expression, for instance:

```
select x from Person x
```

produces an atom bag of objects.

## Struct Atom

The most direct way to construct a struct atom is as follows: struct(\{identifier: expr $\}$ )

For instance:
struct(a: 1)
struct(format: text, $\mathrm{s}:$ "this is the text")
struct(name: "john", age: 32, spouse: first(select Person))

### 5.3 Keywords

Any programming language has its own set of reserved words (keywords) that cannot be used as identifiers. For instance, the keyword "if" cannot be used as a variable in a C program.
OQL also has its own set of keywords. But OQL is one part among others in the information system: for instance, there are an Object Model, an Object Definition Language (ODL) and Language bindings. The Object Model does not introduce any keyword, while ODL has its own set of keywords which are different from the OQL keywords. For instance, a class can include an attribute whose name is "if" as it is not a ODL keyword. If one wants to access this attribute in OQL, using for instance the path expression "x.if", we will get a syntax error. This is not acceptable.
We introduce in OQL (and in ODL) a way to neutralize any keyword: the token "@" used as a prefix keyword neutralizes the keyword and makes it a valid identifier. For instance, "x.@if", denotes the attribute "if" of the instance "x". More generaly, "@identifier" denotes the identifier "identifier" whether "identifier" is a keyword or not.

OQL introduces the following keywords:

| add | all | append | array | as |
| :--- | :--- | :--- | :--- | :--- |
| asc | bag | bodyof | break | by |
| char | classof | contents | define | delete |
| desc | distinct | do | element | else |
| empty | eval | except | exists | false |
| float | for | from | function | group |
| having | ident | if | import | in |
| int | intersect | is | isset | like |
| list | mod | new | nil | not |
| oid | order | pop | print | push |
| refof | return | scopeof | select | set |
| string | struct | structof | suppress | then |
| throw | to | true | typeof | union |
| unset | unval | valof | where | while |

### 5.4 Comments

In OQL, comments are identical to the C++:

- all characters after the token // until the end of the current line are ignored by the interpreter,
- all characters between the tokens /* and */ are ignored.

For instance:

```
1 + 2; // this is a comment
a := "hello"; /* this is another
    comment */
```


### 5.5 Statements

A valid OQL construct is composed of a sequence of statements. Main of the statements are expression statements.

A statement can be one of the following:

- an expression statement,
- a selection statement, if/else,
- an iteration statement, while, do/while, for,
- a function definition statement, function
- a jump statement, break, return
- a compound statement,
- an empty statement.

The OQL expression sub-grammar is very close from the C grammar. The OQL grammar for the flow controls statement - if/else, while, do/while, for - is identical to the C grammar. The common operators of OQL and C have the same associativity and precedence.

### 5.6 Expression Statements

An expression statement is an expression following by a semicolon. Expressions are built from typed operands composed recursively by operators.

The syntax of an expression statement is as follows:
expr ;
where expr denotes any expression.
There are three main kinds of expressions: atomic expressions, unary expressions and binary expressions. Atomic expressions are composed of one terminal atom and no operators, unary expressions are composed of one operand and one operator, binary expressions are composed of two operands and one operator.

We divide the OQL expression family into several semantical sub-families according to their operators as follows:

[^0]In the following sub-sections, we introduce all the OQL expression types using the following template presentation:

1. we present first, in an unformal way, the syntax and the semantics of the operators,
2. general formal information is presented in a first table:

- operator(s)
- type
- syntax
- operand types
- functions

3. in an optionnal second table, we present all the valid operand combination and their result type. A comment about the function performed is added if necessary. This table is skipped in case of the operand type combination is unique or trivial.
4. in a last table, we introduce a few examples. The examples manipulating database objects use the schema that can be found in the directory \$EYEDBROOT/examples/common:
```
// person.odl
enum CivilState {
    Lady = 0x10,
    Sir = 0x20,
    Miss = 0x40
};
class Address {
    attribute string street;
    attribute string<32> town;
    attribute string country;
};
class Person {
    attribute string name;
    attribute int age;
    attribute Address addr;
    attribute Address other_addrs[];
    attribute CivilState cstate;
    attribute Person * spouse inverse Person::spouse;
    attribute set<Car *> cars inverse owner;
    attribute array<Person *> children;
    int change_address(in string street, in string town,
                                    out string oldstreet, out string oldtown);
    static int getPersonCount();
    index on name;
};
class Car {
    attribute string brand;
    attribute int num;
    Person *owner inverse cars;
};
class Employee extends Person {
    attribute long salary;
};
```

Expression types are gathered so to minimize the number of tables in this document.

### 5.7 Atomic Literal Expressions

Atomic literal expressions are expressions composed of a single terminal (or token or lexical unit) without any operator. They are also called primary expressions. These expressions have already been introduced in Section 5.1.

| General Information |  |
| :--- | :--- |
| Operator | no operator |
| Type | unary |
| Syntax | terminal atom |
| Operand Types | integer, float, char, <br> string, boolean, <br> identifier, null, nil, <br> oid, |
| Result Type | same type as the operand |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| 1 | 1 |
| 2. | 12. |
| 'a' | 'a' |
| "hello" | "hello" |
| alpha | value of alpha |
| true | true |
| $83283.1 .29292:$ oid | an error is raised in case of <br> the oid is invalid. <br> Otherwise the re- <br> sult is the input oid: <br> 83283.1.29292:oid |

### 5.8 Arithmetic Expressions

Arithmetic expressions gather the expressions used for any arithmetic computation: addition, multiplication, substraction, division, modulo, shift left and right, bitwise or, bitwise exclusive, bitwise and, bitwise complement. This Section introduced these operators with a special focus on the additive operator which is a multi-purpose operator.

## Additive Expression

The additive operator (i.e. + ) is used for arithmetic addition of integers, floating point numbers and characters, and is also used for string concatenation, list or array concatenation and set or bag union. Its functionality depends on the type of its operands: it is a polymorphic operator. Note that the choice of its functionality is done at evaluation time, not at compile time. That means that the functionality of an expression such as $\mathrm{x}+\mathrm{y}$ is unknown until the evaluation time. Depending on the dynamic type of the operands x and y , it can be a simple arithmetic addition, a string or list or array concatenation, a set or bag union or it can raise an error.
When used as a arithmetic operator and when the two operands have not the same type, one of the operands can be automatically promote to the type of the second one. The promotion mechanism is the same as in the C or $\mathrm{C}++$ languages: integer may be promoted to float, char may be promoted to float or integer.

| General Information |  |
| :--- | :--- |
| Operator | + |
| Type | binary |
| Syntax | expr + expr |
| Commutative | yes |
| Operand Types | integer, float, char, string, list, <br> bag, set, array |
| Result Type | see following table |


| Function | multi functions according to operands: <br> arithmetic addition, string concatenation, <br> list or array concatenation, set or bag <br> union. |
| :--- | :--- |


| Possible Operand Combinations |  |  |  |
| :--- | :--- | :--- | :--- |
| first operand type | second operand type | result type |  |
| integer | integer | integer |  |
| integer | float | float |  |
| char | char | integer |  |
| char | integer | integer |  |
| char | float | float | string concatenation |
| float | float | float | list concatenation |
| string | string | string | array concatenation |
| list | list | list | set union |
| array | array | array | bag union |
| set | set | set |  |
| bag | bag | bag |  |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| $1+2$ | 3 |
| $1+2$. | 3. |
| $2+2.3$ | 4.3 |
| 'a' + 'b' | 195 |
| 'a' +1.2 | 98.200 |
| "hello" + "world" | "helloworld" |
| list (1, 2, 3) + <br> list (2, 3, 4) | list (1, 2, 3, 2, 3, 4) |
| set(1, 2, 3) + <br> set(2, 3, 4) | set(1, 2, 3, 4) |
| $1+$ "hello" | raises an error |
| set (1, 2, 3) + <br> list (2, 3, 4) | raises an error |

## Multiplicative, Division and Minus Expressions

Multiplicative, division and minus expression syntax, semantics, associativity and precedence are quite identical to the corresponding $C$ and $C++$ expressions. When operands have different types, promotionnal mechanisms are the same as for the additive operator.

| General Information |  |
| :--- | :--- |
| Operators | - |
|  | $*$ |
|  | $/$ |
| Type | binary |
| Syntaxes | expr $-\operatorname{expr}$ |
|  | expr * expr |
|  | expr / expr |
| Commutative | $-:$ no |
|  | $*:$ yes |
|  | $/:$ no |
| Operand Types | integer, float, char |
| Result Type | see following table |


| Functions | $-:$ substract |
| :--- | :--- |
|  | *: multiply |
|  | $/:$ divide |


| Possible Operand Combinations |  |  |
| :--- | :--- | :--- |
| first operand type | second operand type | result type |
| integer | integer | integer |
| integer | float | float |
| integer | char | integer |
| char | char | integer |
| char | integer | integer |
| char | float | float |
| float | float | float |
| float | integer | float |
| float | char | float |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| $1-2$ | -1 |
| $3 * 2$. | 6. |
| $2 *{ }^{2}$ 'a' | 194 |
| 'a' * 'b' | 9506 |
| $1 / 2$ | 0 |
| $1 / 2$. | .5000 |
| $1 . / 2$ | .5000 |
| $1 . / 2$ | .5000 |
| "hello" * "world" | raises an error |
| $1-$ "hello" $^{2}$ | raises an error |

Shift, Mod, And, Or, XOr Expressions
Shift, modulo, and, or and xor expression syntax, semantics, associativity and precedence are quite identical to the corresponding C and C++ expressions. Operand types must be integer or char and the only possible type promotion is from char to integer.

| General Information |  |
| :---: | :---: |
| Operators | $\begin{aligned} & \text { << } \\ & \text { >> } \\ & \% \\ & \& \\ & \text { । } \end{aligned}$ |
| Type | binary |
| Syntaxes | $\begin{aligned} & \text { expr << expr } \\ & \text { expr >> expr } \\ & \text { expr \% expr } \\ & \text { expr \& expr } \\ & \text { expr । expr } \\ & \text { expr ~ expr } \end{aligned}$ |
| Commutative | $\begin{array}{ll} \hline \ll & : \\ \gg & \text { no } \\ \% & \text { no } \\ \% & : \\ \text { \&o } \\ \& & : \\ \text { \| yes } & : \end{array}$ |


|  | $\sim:$ yes |
| :--- | :--- |
| Operand Types | integer, char |
| Result Type | integer |
| Functions | $\ll:$ left shift |
|  | $\gg:$ right shift |
|  | $\%$ : modulo |
|  | $\&:$ bitwise and |
|  | $1 \quad:$ bitwise or |
|  | $\sim:$ bitwise exclusive or |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| $1 \ll 4$ | 16 |
| $100 \gg 2$ | 25 |
| $100 \%$ 13 | 9 |
| $0 x f 12$ \& 0xf | 2 |
| $0 \times f 12$ \| 0xf | 3871 |
| $0 \times f 12 ~^{\prime} 0 x f$ | 3869 |
| 'b' \% '9' | 8 |
| $2 \ll 1.2$ | raises an error |
| $2 \% 3.4$ | raises an error |
| $2.1 \% 3$ | raises an error |

## Sign Expressions

Sign expressions are the expressions using the unary operators + or -. The expression syntax, semantics, associativity and precedence are quite identical to the corresponding C and $\mathrm{C}++$ expressions. These unary operators accept only integer, char and float operands.

| General Information |  |
| :--- | :--- |
| Operators | + |
| Type | - |
| Syntaxes | + expr <br> $-e x p r$ |
| Operand Types | integer, char, float |
| Result Type | see following table |
| Functions | sign operator |


| Possible Operand Combinations |  |
| :--- | :--- |
| operand type | result type |
| integer | integer |
| float | float |
| char | integer |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| +12 | 12 |
| -100 | -100 |
| -123.4 | -123.4 |
| '' 'a' $^{\text {-'a' }}$ | 97 |


| +"hello" | raises an error |
| :--- | :--- |
| -null" | raises an error |

## Complement Expressions

The complement operator performs a bitwise complement on its operand. The expression syntax, semantics, associativity and precedence are quite identical to the corresponding $C$ and $C++$ expressions. This operator accepts only integer and char operands.

| General Information |  |
| :--- | :--- |
| Operator | $\sim$ |
| Type | unary |
| Syntax | $\sim$ expr |
| Operand Types | integer, char |
| Result Type | integer |
| Functions | bitwise complement |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| ${ }^{\sim} 112$ | -113 |
| ${ }^{\sim} 0$ | -1 |
| $\sim^{\prime}$ 'a' | -98 |
| $\sim^{\sim} 2.3$ | raises an error |
| $\sim^{\prime}$ "hello" | raises an error |

### 5.9 Assignment Expressions

The expression syntax, semantics, associativity and precedence are quite identical to the corresponding $C$ and $C++$ expressions except that the simple assignment operator in OQL is $:=$ instead of $=$ in C or $\mathrm{C}++$. The left operand must be a left value. A left value is an OQL entity which is assignable: for instance any identifier or a valid path expression.
When the assignment is simple and when the left value is an identifier, no type checking on the right operand is done. For instance, $\mathrm{x}:=10$ and $\mathrm{x}:=$ "hello" are always valid expressions. In the case of the left value is a path expression, the OQL interpreter checks that the type of the second operand matches the expected type of the first one. For instance if $p$ denotes a Person instance, p->age $:=32$ is certainly valid while p->age $:=$ "hello" raises a type check error.
When the assignment is combined with another operation (for instance, the $-=$ operator), the left operand must be initialized and the interpreter checks that the left and right operand can be combined through the other operator.

For instance, the following constructs are valid:

```
a := 10;
a -= 20;
a := "hello";
a += " world";
p := first(select Person);
p.name := "johnny";
first(select Person.age = 0).name := "baby";
```

while these ones produce errors:

```
a := "hello";
a -= 20; // raises the error: operation 'string + integer' is not valid
a := list(1, 2);
a *= 2; // raises the error: operation 'list * integer' is not valid
unset b;
```

b += 20; // raises the error: uninitialized identifier 'b'
p := first(select Person);
p.age := "baby"; // raises the error: integer expected, got string

| General Information |  |
| :---: | :---: |
| Operators | $\begin{aligned} & := \\ & *= \\ & /= \\ & \%= \\ & += \\ & \text { == } \\ & \ll= \\ & \gg= \\ & \&= \\ & \text { l= } \\ & = \\ & = \end{aligned}$ |
| Type | binary |
| Syntaxes | lvalue := expr <br> lvalue $*=$ expr <br> lvalue /= expr <br> lvalue \%= expr <br> lvalue $+=$ expr <br> lvalue -= expr <br> lvalue <<= expr <br> lvalue >>= expr <br> lvalue \&= expr <br> lvalue $\mid=$ expr <br> lvalue ^= expr |
| Commutative | no |
| Operand Types | leftvalue on the left side and any type on the right side |
| Result Type | the type of the right operand |
| Functions | perform an operation and assignment |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| $\mathrm{a}:=24$ | 24 |
| $\mathrm{a}+=12$ | 36 |
| $\mathrm{a} /=2$ | 18 |
| $\mathrm{a}{ }^{\wedge}=100$ | 118 |
| first (select <br> Person) .age $:=$ <br> 38 | 38 |
| "hello" $:=4$ | raises an error (i.e. "hello" is not a left- <br> value) |
| $8:=5$ | raises an error (i.e. 8 is not a leftvalue) |
| unset a; a += 20 faises an error (i.e. uninitialized identi- |  |
| fier) |  |

### 5.10 Auto Increment \& Decrement Expressions

The expression syntax, semantics, associativity and precedence are quite identical to the corresponding $C$ and $C++$ expressions. The operand must be an initialized left value of type integer, char or float. In case of the operand is a char atom, the result type is an integer. Otherwise, the result type is the type of the operand.

| General Information |  |
| :---: | :---: |
| Operators | ++ |
|  | -- |
| Type | unary |
| Syntaxes | expr -- |
|  | expr ++ |
|  | ++expr |
|  | --expr |
| Operand Type | leftvalue of type integer, char or float. |
| Result Type | see following table |
| Functions | expr -- : post-decrementation expr ++ : post-incrementation ++ expr : pre-incrementation --expr : pre-incrementation |


| Possible Operand Combinations |  |
| :--- | :--- |
| operand type | result type |
| integer | integer |
| float | float |
| char | integer |


| Expression Examples |  |  |
| :--- | :--- | :--- |
| expression | result |  |
| $\mathrm{a}:=1 ; \mathrm{a}++$ | 1 | initially a equals 1; the result of the <br> evaluation is 1 but after the evalu- <br> ation, a equals 2 |
| --a | 0 |  |
| $\mathrm{a}++$ | 0 | a equals 1 after the evaluation |

### 5.11 Comparison Expressions

The expression syntax, semantics, associativity and precedence are quite identical to the corresponding C and $\mathrm{C}++$ expressions except that the equal operator could be either $==$ or $=$.

## Equal and NotEqual Expressions

Operands may have any type at all. If the type of the operands differ (modulo the type promotion mechanisms for numbers), the result of the expression operand1 == operand2 is always false while the result of operand1 !=operand2 is always true.

| General Information |  |
| :--- | :--- |
| Operators | $==$ <br> $!=$ |
| Type | binary |
| Syntaxes | expr $==\operatorname{expr}$ <br> expr $!=\operatorname{expr}$ |
| Commutative | yes |
| Operand Types | any type |
| Result Type | boolean |
| Functions | equal <br> not equal |

When operands are number of different types, an automatic promotion is done to the more precise type.

| Expression Examples |  |
| :--- | :--- |
| expression | result |
| $1==1$ | true |
| $1==1.0$ | true |
| $1!=2$ | true |
| $1==2$ | false |
| $1==$ "hello" | false |
| "hello" $==$ "hello" | true |
| list (1, 2, 3) $==$ <br> list | true |
| set $(1, ~ 3, ~ 3) ~$ <br> $\operatorname{set}(1, ~ 2, ~ 3) ~$ | true |

## Less and Greater Expressions

The comparison operators $\langle,\langle=,>$ and $>=$ are multi-purpose operators: they are used for integer, floating point number and character comparison, but also for list or array term-to-term comparison and for set or bag inclusion. Their functionality depends on the type of its operands: they are polymorphic operators. Note that the choice of the functionality is done at evaluation time, not at compile time. That means that the functionality of an expression such as $\mathrm{x}<\mathrm{y}$ is unknown until the evaluation time. Depending on the dynamic type of the operands x and y , it can be an arithmetic comparison (if x and y are numbers), a alpha-numeric comparison (if x and y are strings), a term-to-term ordered collection comparison (if x and y are lists or arrays) or a set or bag inclusion comparison (if x and y are sets or bags).

While arithmetic and alpha-numeric comparisons are trivial and do not need any suplementary explanations, the term-to-term ordered collection comparisons needs to be detailed.

The general algorithm for this functionnality is as follows:

1. let $l 1$ and $l 2$ two OQL ordered collections, containing respectively l1_cnt and l1_cnt atoms.
2. let $o p$ one of the following polymorphic comparison operators: \ll= >>=,
3. $l 1$ op $l 2$ is true if and only if all the following conditions are realized:
(a) $l 1$ and 12 must be of the same collection type,
(b) l1_cnt op l2_cnt or l1_cnt equals l2_cnt
(c) for each atom $l 1[i]$ with $i$ in [i,l1_cnt], l1[i] op $12[i]$

For instance:
$\operatorname{list}(1,2)<=\operatorname{list}(0,2)$ is true
list(1, 2) <= list(3) is false
list (1, 2) <= list(3) is false
list("aaa", 4) < list("bbbb", 8) is true
list("aaa", 4, list(1, 2)) < list("b", 8, list(2, 3)) is true
list (set $(2,4), 3)<\operatorname{list}(\operatorname{set}(4,2,3), 4)$ is true
list (2, 3) < list("hello", 2) raises an error
list $(2,3)<\operatorname{array}(2,4)$ raises an error
Note that the fact that $l 1<=12$ is false does not implie that $l 1>l 2$ is true. Indeed, list(2, 3 ) < list(1, 3 , $2)$ and list(1, 3, 2) >= list(2, 3) are false.

| General Information |  |
| :--- | :--- |
| Operators | $<$ |
|  | $<=$ |
|  | $>$ |
|  | $>=$ |
| Type | binary |
| Syntaxes | $\operatorname{expr}<\operatorname{expr}$ |
|  | $\operatorname{expr}<=\operatorname{expr}$ |
|  | $\operatorname{expr}>\operatorname{expr}$ |
|  | $\operatorname{expr}>=\operatorname{expr}$ |


| Commutative | no |
| :--- | :--- |
| Operand Types | integer, float, char, <br> string, list, array, <br> set, bag |
| Result Type | boolean |
| Functions | the function depends on the <br> operands: <br> < : less than or is included <br> in |
|  | <= : less than or equal or is <br> included in or equal |
|  | $>:$ greater than or contains |
|  | $>=:$ greater than or equal or |
|  | contains or equal |


| Possible Operand Combinations |  |  |  |
| :--- | :--- | :--- | :--- |
| first operand type | second operand type | result type | comments |
| integer, char, float | integer, char, float | boolean | performs an arithmetic comparison |
| string | string | boolean | performs an alpha-numeric compar- <br> ison |
| set | set | boolean | performs an inclusion comparison |
| bag | bag | boolean | performs an inclusion comparison |
| set | bag | boolean | performs an inclusion comparison: <br> the set operand is converted to a <br> bag |
| bag | set | boolean | performs an inclusion comparison: <br> the set operand is converted to a <br> bag |
| list | list | boolean | performs a term-to-term polymor- <br> phic (i.e. numeric, alpha-numeric <br> or inclusion) comparison |
| array | array | boolean | performs a term-to-term polymor- <br> phic comparison |

Note that in case of different operand types, an automatic promotion is done to the more precise type.

| Expression Examples |  |
| :--- | :--- |
| expression | result |
| $1<2$ | true |
| $1>=2$ | false |
| $2 .<=2$ | true |
| "hello" < "world" | true |
| "hello" >= "world" | false |
| list(1, 2) < list(2, 3) | true |
| list(1, 2) < list(0, 3) | false |
| list(1, 2) < list(0, 3, 2) | false |
| list(0, 3, 2) >= list(1, 2) | false |
| list(1, 2) < list(2, 3, 3) | true |
| list(1, 2) < list(0) | false |
| set(1, 2) < set(2, 4, 44) | false |
| set(1, 2) < set(2, 1, 44) | true |
| "hello" >= 3 | raises an error |
| list(1, 2) < array(2, 4, 44) | raises an error |
| set(1, 2) < bag(2, 1, 44) | raises an error |

## Regular Expression Operators

OQL provides the ODMG OQL regular expression operator like plus four other ones. These four extra operators are based on the regular expression UNIX library. So, the syntax of the regular expression are the same as that used by the well known UNIX tools grep, sed, and so on. All the regular expression operators takes two string operands: the first one is the string to compare, the second one is the regular expression. So, these operators are not commutative. These operators provide the following functionalities:
like : ODMG OQL operator. Returns true if the first operand matches the regular expression. Otherwise false is returned. The regular expression is am SQL regular expression where, for instance, \% and - are wilcard characters.
~ : This operator has the same functionnality as the like operator but the regular expression has the UNIX syntax.
~ : Returns true if the first operand matches the regular expression in a case insensitive way. Otherwise false is returned.
!~ : Returns true if the first operand does not match the regular expression. Otherwise false is returned.
!~~ : Returns true if the first operand does not match the regular expression in a case insensitive way. Otherwise false is returned.

Note that the operator like uses currently the UNIX form of regular expressions instead of the SQL form. It will become ODMG/SQL compliant in a next version.

| General Information |  |  |
| :---: | :---: | :---: |
| Operators | $\begin{aligned} & \hline \sim \\ & \sim \sim \\ & !\sim \\ & !\sim \\ & \text { like } \end{aligned}$ |  |
| Type | binary |  |
| Syntaxes | expr ~ expr <br> expr ~~ expr <br> expr !~ expr <br> expr ! ~~ expr <br> expr like expr |  |
| Commutative | no |  |
| Operand Types | string |  |
| Result Type | boolean |  |
| Functions | !~: !~~ : <br> like: | matches the regular expression matches the regular expression, case insensitive does not match the regular expression does not match the regular expression, case insensitive matches the regular expression |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| "hello" ~ "LL" | false |
| "hello" ~~ "LL" | true |
| "hello" ~ "^LL" | false |
| "hello" ~ "^h" | true |
| "hello" !~ "^h" | false |
| "hello" ~ ".*ll.*" | true |
| ".*ll.*" ~ "hello" | false because regular ex- <br> pression should be on the <br> right |

### 5.12 Logical Expressions

OQL provide three logical expressions which can take two form each. The logical or operator is \| or or. The logical and operator is \&\& or and. The logical not operator is ! or not.

The expression syntax, semantics, associativity and precedence are quite identical to the corresponding C and $\mathrm{C}++$ expressions. Note that the ODMG operator "II" denotes the string concatenation.
As for C and C++, the OQL interpreter performs a lazy evaluation:

- expr1 || expr2 expr2 is not evaluated if expr1 is evaluated to true.
- expr1 \&\& expr2
expr2 is not evaluated if expr1 is evaluated to false.

| General Information |  |
| :--- | :--- |
| Operators | II <br>  |
| Type | unary, binary |
| Syntaxes | expr II expr <br> expr or expr <br> expr \&\& expr <br> expr and expr <br> $!$ expr <br> not expr |
| Operand Type | boolean |
| Result Type | boolean |
| Functions | logical or <br> logical and <br> logical not |


| Expression Examples |  |
| :---: | :---: |
| expression | result |
| true \\| false | true |
| false \|| false | false |
| true \&\& false | false |
| 1 == 2 \|| 3 == 4 | false |
| 1 == 2 \|| $3==3$ | true |
| $1==2$ or $3==3$ | true |
| $\begin{aligned} & 1==2 \text { I\| "hello" } \\ & ==\text { "hello" } \end{aligned}$ | true |
| $\begin{aligned} & (1==2 \mid 12==2) \\ & \& \&(a=\text { hello" }) \end{aligned}$ | returns true if a equals "hello". false otherwise |
| 1 \|| 3 == 3 | raises an error: boolean expected got integer |
| !3 | raises an error: boolean expected got integer |
| ! (1 == 1) | false |
| $\operatorname{not}(1==1)$ | false |

### 5.13 Conditional Expression

The unique conditional expression operator is ?:. The expression syntax, semantics, associativity and precedence are quite identical to the corresponding C and C++ expressions. The first operand must be an boolean and the two others may be of any type. Contrary to C and C++, the two last operands does not need to be of the same type.

| Operator | $?:$ |
| :--- | :--- |
| Type | ternary |
| Syntaxe | expr ? expr : expr |
| Operand Types | first operand is boolean, oth- <br> ers are any type |
| Result Type | type of the evaluated <br> operand |
| Functions | conditional evaluation: eval- <br> uates and returns the second <br> operand if the first operand <br> is true; otherwise evalu- <br> ates and returns the second <br> operand |


| Expression Examples |  |  |  |
| :--- | :--- | :---: | :---: |
| expression result |  |  |  |
| true ? "hello" : "world" | "hello" |  |  |
| true ? 2.3 : "world" | 2.3 |  |  |
| $1+1:=2$ ? (a := 3.1415926535) : nil | 3.1515926535 |  |  |
| $1 ? 3:$ nil anses an error: |  |  |  |

### 5.14 Expression Sequences

The expression sequence operator - also called comma sequencing - expression syntax, semantics, associativity and precedence are quite identical to the corresponding $C$ and $C++$ expressions. This operator , takes two operands: it evaluates both of them and returns the second one.

| General Information |  |
| :--- | :--- |
| Operator | , |
| Type | binary |
| Syntaxe | expr , expr |
| Commutative | no |
| Operand Types | any type |
| Result Type | type of the second operand |
| Functions | evaluates the first operand, then the sec- <br> ond one. Returns the evaluation of the <br> second one. |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| true, "hello" | "hello" |
| $\mathrm{a}:=2,4$ | 4 (note that a equals 2) |
| $\mathrm{b}:=10, \mathrm{a}:=\mathrm{b}+1$ | 11 |

### 5.15 Array Deferencing

OQL provides polymorphic single and range deferencing. The single deferencing is used to get one element in an ordered collection or a character in a string or one element in a non-collection array. The range deferencing is used to get several elements.
The deferencing of ordered collections is introduced in more details in Section 5.24.

## Single Deferencing

The single deferencing operator takes as its first operand an atom of type string, an indexed (or ordered) collection (list or array) or a non-collection array. The second operand must be of type integer. Depending on the type of the first operand, the returned atom is as follows:

1. if the first operand is a string, the returned atom is the \#expr character of the string where expr denotes the second operand. If expr is equal to the length to the string the character ' $\backslash 000$ ' is returned. If it greater than the length the string an out of bounds error is raised.
2. for an ordered collection, the returned atom is the \#expr item of the collection. If expr is greater than or is equal to the size of the collection, an out of bounds error is raised.
3. if the first operand is an non-collection array, the returned atom is the \#expr item of the array. If expr is greater than or is equal to the size of the array, an out of bounds error is raised.

The single deferencing operator may be used as a left value, that means that a single deferencing expression is assignable. For instance the sequence of statements:

```
s := "hello";
s[1] := 'E';
s[4] := '0';
```

set the variable s to "hEllO".
The single deferencing operator may be used everywhere in a path expression. For instance, first (select Person). other_addrs [2] .str denotes the character \#3 of the street attribute in the \#2 other_addrs non-collection array attribute of the first Person instance.

| General Information |  |
| :--- | :--- |
| Operator | [] |
| Syntaxe | expr $[$ expr $]$ |
| Type | binary |
| Commutative | no |
| Operand Types | first operand: string, indexed collection (list or <br> array) or non-collection array, second operand: <br> integer |
| Result Type | char if first operand is a string, otherwise type <br> of the returned item in the indexed collection or <br> non-collection array. |
| Functions | [expr $]:$ returns the character (or item in the <br> indexed collection or in the non-collection array) <br> number expr |
| Note | this operator may be used in the composition of <br> a left value. |


| Expression Examples |  |  |
| :---: | :---: | :---: |
| expression | result |  |
| "hello" [0] | 'h' |  |
| a := "hello"; a[1] | 'e' |  |
| a[3] | 1 |  |
| a[6] | raises an error |  |
| a[0] : $=$ ' $\mathrm{H}^{\prime}$ | 'H' | a equals "Hello" |
| list(1, 2, "hello", 4) [3] | "hello" |  |
| list(1, 2, "hello", 4) [4] | raises an er- ror |  |
| $\begin{aligned} & \text { first(select } \\ & \text { Person). name [2] := 'X, } \end{aligned}$ | ' X ' |  |

## Range Deferencing

The range deferencing operators, [:] and [?], takes as their first operand an atom of type string, an indexed (or ordered) collection (list or array) or a non-collection array. The other operands must be of type integer. The [?] may have also an unordered collection (set or bag as its first operand.

The operator syntax and semantics are as follows:

- expr [expr1: expr2]

1. if the first operand is a string, the returned atom is a list composed of the characters between the \#expr1 and the \#expr2 characters of the string. If expr1 is less than zero or or if expr2 is greater than the length of the string an out of bounds error is raised.
2. for an ordered collection, the returned atom is a list composed of the items between the \#expr1 and the \#expr2 items of the collection. If expr1 is less than zero or if i expr2 iss greater than or is equal to the size of the collection, an out of bounds error is raised.
3. if the first operand is an non-collection array, the returned atom is a list composed of the items between the \#expr1 and the \#expr2 items of the array. If expr1 is less than zero or if i expr2 iss greater than or is equal to the size of the collection, an out of bounds error is raised.

- expr [?]

1. if the first operand is a string, the returned atom is a list composed of all the characters of the string, including the last character ' $\backslash 000$ '.
2. for an ordered or an unordered collection, the returned atom is a list composed of all the items of the collection. If the collection is a list, the list itself is returned. If the collection is an array, this operator has the same functionnality as the listtoarray library function.
3. if the first operand is an non-collection array, the returned atom is a list composed of the all the items of the array.

Contrary to the single deferencing operator, the range deferencing operator cannot be used as a left value.
The range deferencing operators may be used everywhere in a path expression. For instance, first (select Person). children [?] .name [ denotes the list of all characters of the name attribute in all the children of the first Person instance.

| General Information |  |
| :---: | :---: |
| Operators | [:] |
|  | [?] |
| Syntaxes | $\begin{aligned} & \operatorname{expr}[\operatorname{expr}: \operatorname{expr}] \\ & \operatorname{expr}[?] \end{aligned}$ |
| Type | ternary or unary |
| Operand Types | first operand: string or indexed collections (list or array), second operand and third operand: integer |
| Result Type | a list of char if first operand is a string, otherwise a list of returned items in the indexed collection or non-collection array. |
| Functions | [expr1: expr2] : returns a lits of characters (or items in collection) indexed from expr 1 to expr2 [?] : returns a list of all characters (or items in collection) |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| "hello" [0:2] | list('h', 'e', 'l') |
| "hello" [?] | list('h', 'e', 'l', 'l', 'o', <br> ' $\backslash 000$ ') |
| list(1, 2, "hello", 4) [2:3] | list("hello", 4) |
| array(1, 2, "hello", 4) [?] | list(1, 2, "hello", 4) |


| first(select Person).name[?] | list('j', 'o', 'h', 'n', <br> $, \backslash 000 ')$ |
| :--- | :--- |
| list(select class.type $=$ <br> "user") [0:4].name | list("Employee", "Address", <br> "Person") |

### 5.16 Identifier Expressions

We call an identifier expression an unary or binary expression whose operands must be identifiers. There are height identifier operators: : : isset, unset, \& (identical to refof), * (identical to valof), scopeof, push and pop.

As all these operators take identifiers as their operands, we skip the second table (operand combinations) while introducing these operators.

## : : Operator

The : : unary/binary operator (called scope operator) is used to define a global or particular scope for a variable. The unary version of this operator denotes a global scope. For instance, : : alpha denotes the global variable alpha. In the body of a function, identifiers denote local variables; outside the body of a function identifiers denote global variables, that means that, in this context, the global scope operator is not mandatory. If one wants to use a global variable in the body of a function, the global scope operator is mandatory. Refer to the Function Definition Statement Section for more information about local function variables.

The binary version of this operator denotes a particular scope. For instance, Person: :checkName denotes the class attribute or method of the class Person.
note: class (or static) attributes are not currently well supported by the OQL interpreter. Class attributes are only supported in some specific query expressions (refer to the Query Expression Section).

| General Information |  |
| :--- | :--- |
| Operator | $::$ |
| Syntax | $::$ identifier <br> identifier: :identifier |
| Type | unary and binary |
| Operand Types | identifier |
| Result Type | value of the identifier if used as a right value; iden- <br> tifier reference if used as a left value |
| Function | defines a global or particular scope for the identi- <br> fier. |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| $:: \mathrm{a}$ | the value of the global vari- <br> able a |
| $::$ alpha $:=1$ | sets the value of the global <br> variable alpha to 1, returns <br> 1 |
| Person: :checkName("wayne") | calls the class method <br> checkName in the class <br> Person |
| 2::alpha | raises an error |

## isset Operator

The isset operator is used to check whether a variable is already set or not. It returns true is the variable is set, false otherwise.

| Operator | isset |
| :--- | :--- |
| Syntax | isset identifier |
| Type | unary |
| Operand Type | identifier |
| Result Type | boolean |
| Function | evaluated to true if the identifier is set, false <br> otherwise |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| isset oql\$variables | true |
| isset a | returns true if a is set, <br> false otherwise |
| isset 1 | raises an error |

## unset Operator

The unset operator is used to unset an variable. It returns the nil atom.

| General Information |  |
| :--- | :--- |
| Operator | unset |
| Syntax | unset identifier |
| Type | unary |
| Operand Type | identifier |
| Result Type | nil |
| Function | unset the identifier |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| unset a | nil |
| unset : :a | nil |
| unset 2 | raises an error |

## refof Operator

The \& (identical to refof) operator is used to get the reference of an identifier. This operator is essentially used when one calls a function or method which updates one or more given parameters. For instance, let the function swap ( $x$, $y$ ) which swaps the value of its two parameters. One needs to give the references of the variables that one wants to swap. For instance:

```
i := "ii";
j := "jj";
```

swap(\&i, \&j);
After the call to swap, the variable i equals jj while the variable $j$ equals ii.
The reverse operator $*$ (described following section) is used in the swap function.

| General Information |  |
| :--- | :--- |
| Operator | refof <br>  |
| Syntax | refof identifier <br> \& identifier |
| Type | unary |


| Operand Type | identifier |
| :--- | :--- |
| Result Type | identifier |
| Function | evaluates the expression to the <br> identifier reference; returned an <br> identifier atom |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| \&alpha | alpha |
| refof alpha | alpha |

## valof Operator

The * (identical to valof) operator is used to get the value of the identifier pointed by a reference. For instance, after the two following expressions:
alpha := 1;
ralpha := \α
*ralpha equals 1.

This operator may be used in the composition of a left value, for instance:

```
alpha := 1;
ralpha := &alpha;
*ralpha := 2; // now, alpha equals 2
*ralpha += 8; // now, alpha equals 10
```

But this operator is essentially used in the body of functions or methods which update one or more given parameters, for instance, the function swap described in the previous section is as follows:

```
function swap(x, y) {
        v := *x;
    *x := *y;
    *y := v;
}
```

| General Information |  |
| :--- | :--- |
| Operator | valof <br> $*$ |
| Syntax | valof identifier <br> $*$ identifier |
| Type | unary |
| Operand Type | identifier |
| Result Type | value of the identifier |
| Function | returns the value of the identifier <br> denotes by the operand |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| *alpha | if alpha value is an atom <br> identifier $x$, returns the value <br> of $x$, otherwise an error is <br> thrown |
| $\mathrm{x}:=12$; alpha $:=\& \mathrm{x} ; * \mathrm{x}$ | $* \mathrm{x}$ returns 12 |

## scopeof Operator

The scopeof operator returns the string "global" or "local" depending whether the identifier is global or local.

| General Information |  |
| :--- | :--- |
| Operator | scopeof |
| Syntax | scopeof identifier |
| Type | unary |
| Operand Type | identifier |
| Result Type | string |
| Function | returns the scope of the identifier. |


| Expression Examples |  |  |
| :--- | :--- | ---: |
| expression | result |  |
| scopeof ::alpha | returns "global" for any <br> alpha if it set; otherwise an <br> error is thrown |  |
| scopeof alpha | returns "global" or <br> "local" depending on <br> the context. |  |

## push Operator

The push operator is used to push an identifier on a new local table. This operator is rarely used.

| General Information |  |
| :--- | :--- |
| Operator | push |
| Syntax | push identifier <br> push identifier $:=$ expr |
| Type | unary and binary |
| Operand Types | first operand identifier, optionnal second <br> operand: any type |
| Result Type | identifier or any type in case of an assignment |
| Function | push the identifier on to the symbol table stack. <br> An assignment can be performed at the same <br> time. Returns the identifier or the value of the <br> expression assignment. |


| Expression Examples |  |
| :--- | :--- |
| expression | result |
| push a | pushes a on a new local sym- <br> bol table. |
| push a $:=10$ | pushes a on a new local <br> symbol table and assigns its <br> value to 10 |

## pop Operator

The pop operator is used to pop an identifier from a local table. It is used after a push. For instance:

```
a := "hello";
a; // a equals "hello"
push a := 10;
a;
pop a;
a; // a equals "hello"
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & pop \\
\hline Syntax & pop identifier \\
\hline Type & unary \\
\hline Operand Type & identifier \\
\hline Result Type & the type of the value of the identifier \\
\hline Function & pop the identifier from the symbol table stack \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline expression & result \\
\hline pop a & \begin{tabular}{l} 
returns the value of a if it \\
is set; otherwise an error is \\
returned
\end{tabular} \\
\hline
\end{tabular}

\subsection*{5.17 Path Expressions}

The path expression operator \(\rightarrow\) (identical to .) is used to navigate from an object and read the right data one needs. This operator enables us to go inside complex objects, as well as to follow simple relationships. For instance, if p denotes a Person instance, p.spouse denotes the spouse attribute of this person.
The more complex expression p.spouse.address.street denotes the street in the address of spouse of the person \(p\). This notation is very intuitive because it looks like the well known C, C++ and Java syntaxes.

The path expression operator may composed a left value, for instance:
p.spouse.name := "mary";
set the name of the spouse of the person \(p\) to mary.

This operator may be combined with the array deferencing operators, for instance:
p.spouse. name [2];
p.spouse.name[2] := 'A';
p.spouse.other_addrs[2].street[3] := 'C';
p.spouse.children[?];
p.spouse.children [?]. name;

The path expression operator may be also used to navigate through struct atom, for instance: (struct (a : 1, b : "hello")).b returns "hello". Note that because of the precedence of operators, parenthesis are necessary around the literal struct construct.
Finally, the path operator may be applied to a collection; in this case a collection of the same type of this operand is returned. For instance:
(select Person). name returns a bag of string.
(select distinct Person). age returns a set of int.
Note that the path expression operator is used frequently in the query expressions as shown in a next section.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operators & \begin{tabular}{l}
\(->\) \\
\(->\)
\end{tabular} \\
\hline Syntaxes & \begin{tabular}{l} 
expr \(\cdot \operatorname{expr}\) \\
expr -> expr
\end{tabular} \\
\hline Type & binary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: oid or object, second operand: \\
identifier
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline Result Type & \begin{tabular}{l} 
type of the attribute denoted by the second \\
operand
\end{tabular} \\
\hline Functions & \begin{tabular}{l} 
returns the attribute value denoted by second \\
operand of the object denoted by the first operand \\
The first operand must denote an EyEDB in- \\
stance (object or literal) of an agregat including \\
the attribute denoted by the second operand.
\end{tabular} \\
\hline Note & these two operators are identical \\
\hline
\end{tabular}
\begin{tabular}{|l|l|l|}
\hline \multicolumn{3}{|c|}{ Expression Examples } \\
\hline expression & result & comments \\
\hline p ->name & \begin{tabular}{l} 
the value of attribute \\
name in the object de- \\
noted by p
\end{tabular} & \begin{tabular}{l}
p must denote an \\
EYEDB instance (ob- \\
ject or literal) of an \\
agregat including the \\
attribute name
\end{tabular} \\
\hline \begin{tabular}{l} 
first (select \(\times\) Person \\
x from \(\mathrm{x} . l\) lastname \(=\) \\
"wayne")->lastname
\end{tabular} & "wayne" & \\
\hline
\end{tabular}

\subsection*{5.18 Function Call}

OQL allows one to call an OQL function with or without parameters. The operator for function call is (). A function call may be the first term of a path expression, for instance: first(select Person)->name.
Contrary to the method invocation, there are no function overloading mechanisms: that means, that one cannot have differents functions with the same name and a different signature. To take benefit of the overloading mechanisms, one must use methods.
Note: contrary to the ODMG 3 specifications, one currently needs to use parenthesis to invoke a method even if the method has no arguments.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & () \\
\hline Syntaxe & expr (expr_list) \\
\hline Type & n-ary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: identifier, other operands: any \\
type
\end{tabular} \\
\hline Returned type & type of the returned atom by the function call \\
\hline Functions & \begin{tabular}{l} 
calls the OQL function denoted by the first \\
operand using the other operands as arguments. \\
The number of operands must be equal to the \\
number of arguments of the OQL function plus \\
one
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline expression & result \\
\hline fact(10) & 3628800 \\
\hline fact(fact(3)) & 720 \\
\hline toUpper("hello world") & "HELLO WORLD" \\
\hline toUpper("hello") + "world" & "HELLOworld" \\
\hline interval(1, 5) & \begin{tabular}{l} 
list(1, 2, 3, 4, \\
5)
\end{tabular} \\
\hline swap(\&i, \&j) & nil \\
\hline \begin{tabular}{l} 
first(select \\
Person).spouse.name
\end{tabular} & "mary" \\
\hline
\end{tabular}

\subsection*{5.19 Method Invocation}

\section*{Instance Method Invocation}

OQL allows one to call a instance method with or without parameters. The method can be written in C++ or in OQL. As in \(\mathrm{C}++\), method calls use a combination of the path expression operator and the function call operator.

As in C++ or Java, methods can be overloaded: that means that one can have differents methods with the same name and a different signature or differents methods with the same name and the same signature in a class hierarchy. The choice of the method to invoke is done at evaluation time not at compile time. For instance let two methods Person Person:: \(f\) (in int, in int) and int Person: :f(in float, in string), the method to be invoked in the expression \(p->f(x, y)\) is decided at evaluation time according to the true types of \(x\) and \(y\) :
```

p := first(select Person);
x := 1; y := 2;
p->f(x, y); // X::f(in int, in int) is invoked
p->f(x, y)->name; // this is valid because p->f(x, y) returns a Person
x := 1.3; y := "hello";
p->f(x, y); // X::f(in float, in string) is invoked
p->f(x, y)->name; // this is not valid because p->f(x, y) returns an integer

```

A major contribution of object orientation is the possibility of manipulating polymorphic objects and thanks to the late binding mechanism to carry out generic actions on the elements of these objects.
For instance, let the two methods void Person::doit(in int) and void Employee::doit(in int), the method to be invoked in the expression p ->doit ( x ) is decided at evaluation time according to the true type of p :
```

p := new Person();
p->doit(1); // Person::doit(in int) is invoked
p := new Employee();
p->doit(1); // Employee::doit(in int) is invoked

```

To invoke a method, the following conditions must be realize:
1. the object or oid on which the method is applied must be an instance of a class, for instance X .
2. the name of the invoked method, the number and the type of arguments must be compatible with an existing method in the class X ,
3. the result type must match the expected type in the expression.

For instance, let the methods int compute(in int, int float) and int compute(in int, in float, in int [], out string) in the class \(X\). To invoke the first method on an instance of \(X\), one needs to apply the method compute to an instance of X with one integer and one float, for instance:
```

x := new X();

```
x.compute (1, 2.3);
x.compute(a := fact(10), float(fib(10)));

To invoke the second method on an instance of X , one needs to apply the method compute to an instance of X with an integer, a float, an ordered collection of integer and a reference to a variable, for instance:
```

x.compute(1, 23.4, list(1, 2, 3, 4), \&a);

```

The following table shows the mapping (which defines the compatibility) between the ODL and the OQL types.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{\(O D L / O Q L\) Mapping } \\
\hline ODL Type & OQL Type \\
\hline \hline in int16 & integer \\
\hline out int16 & identifier \\
\hline inout int16 & identifier initialized to an integer \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline in int32 & integer \\
\hline out int32 & identifier \\
\hline inout int32 & identifier initialized to an integer \\
\hline in int64 & integer \\
\hline out int64 & identifier \\
\hline inout int64 & identifier initialized to an integer \\
\hline in byte & char \\
\hline out byte & identifier \\
\hline inout byte & identifier initialized to a char \\
\hline in char & char \\
\hline out char & identifier \\
\hline inout char & identifier initialized to a char \\
\hline in string & string \\
\hline out string & identifier \\
\hline inout string & identifier initialized to a string \\
\hline in float & float \\
\hline out float & identifier \\
\hline inout float & identifier initialized to a float \\
\hline in oid & oid \\
\hline out oid & identifier \\
\hline inout oid & identifier initialized to a oid \\
\hline in object * & oid of any class \\
\hline out object * & identifier \\
\hline inout object * & identifier initialized to an oid of any class \\
\hline in X * ( X denotes a class instance) & oid of class X \\
\hline out \(\mathrm{X} *\) ( X denotes a class instance) & identifier \\
\hline inout X * (X denotes a class instance) & identifier initialized to a oid of class X \\
\hline in X * [] ( X denotes a class instance) & ordered collection of oid of class X \\
\hline out X * [] (X denotes a class instance) & identifier \\
\hline inout X *[] (X denotes a class instance) & identifier initialized to an ordered collection of oid of class X \\
\hline in X [] (X denotes any ODL type) & ordered collection of atoms bound to X \\
\hline out X [] (X denotes any ODL type) & identifier \\
\hline inout X [] (X denotes any ODL type) & identifier initialized to an ordered collection of atoms bound to X \\
\hline
\end{tabular}

Note: contrary to the ODMG 3 specifications, one currently needs to use parenthesis to invoke a method even if the method has no arguments.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operators & \begin{tabular}{l}
\(->\)
\end{tabular} \\
\hline Syntaxes & \begin{tabular}{l} 
expr . expr (expr_list) \\
expr \(\rightarrow\) expr (expr_list)
\end{tabular} \\
\hline Type & n-ary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: oid or object, second operand: \\
identifier, other operands: any type
\end{tabular} \\
\hline Result Type & type of the atom returned by the method call \\
\hline Functions & \begin{tabular}{l} 
invokes the method denoted by the second \\
operand applied to the object denoted by the first \\
operand, using the other operands as arguments.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|} 
& \begin{tabular}{l} 
The first operand must denote an EYEDB in- \\
stance (object or literal) of an agregat including \\
the method whose name is the second operand. \\
The number of arguments and the type of argu- \\
ments must match one of the methods included \\
in the class of the object denoted by the first \\
operand.
\end{tabular} \\
\hline Note & these two operators are identical \\
\hline
\end{tabular}
\begin{tabular}{|l|l|l|}
\hline \multicolumn{3}{|c|}{ Expression Examples } \\
\hline expression & result & \begin{tabular}{l} 
comments \\
object denoted by p \\
p ->getOid()
\end{tabular} \\
\hline img->compute(1, 2.3) & \begin{tabular}{l} 
the value returned by \\
the method call \\
tive method of the \\
class object, each \\
object can call this \\
method
\end{tabular} \\
\hline \begin{tabular}{l} 
first (select Person.name first operand \\
must denote an \\
"wayne").getSpouse()
\end{tabular} & \begin{tabular}{l} 
the value returned by \\
EyEDB instance \\
the method call
\end{tabular} & \begin{tabular}{l} 
object or literal) of \\
an agregat including \\
the method whose \\
name is compute
\end{tabular} \\
\hline
\end{tabular}

\section*{Class Method Invocation}

OQL allows one to call a class method with or without parameters. The method can be written in C++ or in OQL. As in \(\mathrm{C}++\), method calls use a combination of the scope operator and the function call operator. To invoke a class method, the following conditions must be realize:
1. the name of the invoked method, the number and the type of arguments must be compatible with an existing method in the class X ,
2. the result type must match the expected type in the expression.

The overloading and the late binding mechanisms are the same as for the instance method invocations.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & \(::\) \\
\hline Syntaxe & identifier: :identifier(expr_list) \\
\hline Type & n-ary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: identifier, second operand: \\
identifier, other operands: any type
\end{tabular} \\
\hline Result Type & type of the atom returned by the method call \\
\hline Functions & \begin{tabular}{l} 
invokes the class method denoted by the second \\
operand applied to the class denoted by the first \\
operand, using the other operands as arguments. \\
The first operand must denote an EYEDB class of \\
an agregat including the class method whose name \\
is the second operand. The number of arguments \\
and the type of arguments must match one of the \\
class methods included in the class denoted by the \\
first operand.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline expression & result & comments \\
\hline EyeDB: :getVersion() & 2.8 .8 & \begin{tabular}{l} 
getVersion() is a \\
native static method \\
of the class EyeDB
\end{tabular} \\
\hline Person: checkName("johnny") & \begin{tabular}{l} 
the value returned by \\
the method call
\end{tabular} & \begin{tabular}{l} 
the class method \\
checkName must \\
exist in the class \\
Person and must \\
take one and only \\
one input string \\
argument.
\end{tabular} \\
\hline
\end{tabular}

\subsection*{5.20 Eval/Unval Operators}

\section*{eval Operator}

One major feature of OQL is that one can invoke its evaluator using the eval operator. This allows us to build OQL constructs at runtime and perform their evaluation. This is very useful, for instance, when we want to build a query expression where the projection or the from reference sets are is unknown. For instance, the following function allows us to retrieve the values of the attribute attrname in the class classname:
```

function getValues(classname, attrname) {
cmd := "select x." + attrname + " from " + classname + " x";
return (eval cmd);
}

```
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & eval \\
\hline Syntaxe & eval string \\
\hline Type & unary \\
\hline Operand Types & string \\
\hline Functions & \begin{tabular}{l} 
calls the OQL evaluator on the string operand. \\
The string operand can contain any OQL valid \\
construct: an expression, a statement or a se- \\
quence of statements.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline expression & result \\
\hline eval "10" & 10 \\
\hline eval "a \(:=100\) " & \begin{tabular}{l} 
result is 100; the variable a is set to \\
100
\end{tabular} \\
\hline eval "a := \"hello\"; b := a + \"world\"" & \begin{tabular}{l} 
result is "hello world"; the vari- \\
able a is set to "hello"; the variable \\
b is set to "hello world"
\end{tabular} \\
\hline
\end{tabular}

\section*{unval Operator}

The unval is the inverse of the unval in the sense that it takes any valid OQL expression and returns the string representation; the comments and, when not necessary, the spaces and tabulations are skipped. For instance, the construct unval \(\mathrm{a}:=10\) returns " \((\mathrm{a}:=10)\) ".
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & unval \\
\hline Syntax & unval expr \\
\hline Type & unary \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline Operand Types & any type \\
\hline Functions & returns the string expression \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline expression & result \\
\hline unval 10 & "10" \\
\hline unval alpha += 10 - beta + 1 & "(alpha:=(alpha+((10-beta)+1))) \\
\hline eval unval alpha := "hello" & \begin{tabular}{l} 
returns "hello"; alpha is set to \\
"hello"
\end{tabular} \\
\hline
\end{tabular}

\subsection*{5.21 Set Expressions}

OQL allows us to perform the following operations on sets and bags: union, intersection, difference and inclusion. The operands can be sets or bags. For all these operators, when the operand's collection types are different (bag and set), the set is first converted to a bag and the result is a bag.

\section*{union Operator}

The union operator performs the union of two sets or bags. This operator has the same precedence as the logical or operator.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & union \\
\hline Syntax & expr union expr \\
\hline Type & binary \\
\hline Operand Types & set or bag \\
\hline Result Type & \begin{tabular}{l} 
set if both two operands are of type set, bag \\
otherwise
\end{tabular} \\
\hline Functions & returns the union of the two operands. \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline expression & result \\
\hline \(\operatorname{set}(1,2)\) union \(\operatorname{set}(2,3)\) & \(\operatorname{set}(1,2,3)\) \\
\hline \(\operatorname{set}(1,2)\) union \(\operatorname{bag}(2,3)\) & \(\operatorname{bag}(1,2,2,3)\) \\
\hline list \((1,2)\) union \(\operatorname{bag}(2,3)\) & raises an error \\
\hline
\end{tabular}

\section*{intersect Operator}

The intersect operator performs the intersection of two sets or bags. This operator has the same precedence as the logical and operator.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & intersect \\
\hline Syntax & expr intersect expr \\
\hline Type & binary \\
\hline Operand Types & set or bag \\
\hline Result Type & \begin{tabular}{l} 
set if both two operands are of type set, bag \\
otherwise
\end{tabular} \\
\hline Functions & returns the intersection of the two operands. \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline expression & result \\
\hline \(\operatorname{set}(1,2)\) intersect \(\operatorname{set}(2,3)\) & \(\operatorname{set}(2)\) \\
\hline \(\operatorname{set}(1,2)\) intersect \(\operatorname{bag}(2,3)\) & \(\operatorname{bag}(2)\) \\
\hline \(\operatorname{bag}(1,2,2,3)\) intersect \(\operatorname{bag}(2,3,2)\) & \(\operatorname{bag}(2,2,3)\) \\
\hline list 1,2\()\) intersect \(\operatorname{bag}(2,3)\) & raises an error \\
\hline
\end{tabular}

\section*{except Operator}

The except operator performs the difference between two sets or bags. This operator has the same precedence as the logical or operator.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & except \\
\hline Syntax & expr except expr \\
\hline Type & binary \\
\hline Operand Types & set or bag \\
\hline Result Type & \begin{tabular}{l} 
set if both two operands are of type set, bag \\
otherwise
\end{tabular} \\
\hline Functions & returns the difference of the two operands. \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline expression & result \\
\hline set \((1,2)\) except \(\operatorname{set}(2,3)\) & \(\operatorname{set}(1)\) \\
\hline \(\operatorname{set}(1,2)\) except \(\operatorname{bag}(2,3)\) & \(\operatorname{bag}(1)\) \\
\hline \(\operatorname{set}(1,2,10)\) except \(\operatorname{bag}(12)\) & \(\operatorname{bag}(1,2,10)\) \\
\hline list \((1,2)\) except \(\operatorname{bag}(2,3)\) & raises an error \\
\hline
\end{tabular}

\section*{Inclusion Operators}

The inclusion operators for sets and bags are the comparison operators less than/greater than introduced in a previous section.
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|r|}{General Information} \\
\hline Operator & \[
\begin{aligned}
& \text { < } \\
& <= \\
& > \\
& >=
\end{aligned}
\] \\
\hline Syntax & \[
\begin{aligned}
& \text { expr < expr } \\
& \text { expr <= expr } \\
& \text { expr > expr } \\
& \text { expr >= expr }
\end{aligned}
\] \\
\hline Type & binary \\
\hline Operand Types & set or bag \\
\hline Result Type & boolean \\
\hline Functions & coll1 < coll2 : returns true if and only if coll1 is included in coll2 but not equal to coll2 coll1 > coll2 : returns true if and only if coll2 is included in coll1 and not equal to coll1 coll2 <= coll1 : returns true if and only if coll1 is included in coll2 or equal to coll2 coll1 >= coll2 : returns true if and only if coll2 is included in coll1 or equal to coll1 \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline expression & result \\
\hline \(\operatorname{set}(1,2)<\operatorname{set}(2,3)\) & false \\
\hline \(\operatorname{set}(1,2)<\operatorname{set}(2,3,1)\) & true \\
\hline \(\operatorname{set}(1,2)<\operatorname{bag}(2,3,1)\) & true \\
\hline \(\operatorname{set}(1,2)<=\operatorname{bag}(2,1)\) & false \\
\hline \(\operatorname{set}(1,2)>=\operatorname{bag}(2,1)\) & false \\
\hline
\end{tabular}

\subsection*{5.22 Object Creation}

OQL allows us to create persistent or transient objects using the new operator. The general syntax for an object creation is as follows:
[new] [<[expr]>] class_name( \(\{\) path_expression : expr \(\}\) )
1. the operator new is optionnal: when the operator is missing, the construct is called an implicit new construct. When the optionnal following construct "<[expr]>" is not used, there is no functionnal differences between using new or not.
2. the optionnal construct after the new operator indicates the target location of the object to create:
(a) if this construct is omitted, the object will be a persistent object created in the current default database,
(b) if this construct is under the form <expr>, the OQL interpreter expects for a database object handle as the result of the expression evaluation. This database will be used as the target location. For instance:
```

new < oql\$db > Person();

```
will create a Person instance in the database pointed by oql\$db, which is in fact the current database.
(c) if this construct is under the form <>, the object will be a transient object.
3. the class_name indicates the name of a valid user type in the context of the current database.
4. the path_expression indicates an attribute name or a sequence of attributes using the optional array operator, for instance the following path expressions are valid for an object construction:
```

name
lastname
addr.street
addr.town[3]
spouse.name

```
5. the expr behind path_expression is any OQL expression as soon as its result type matches the expected type of the path_expression.
6. the order of evaluation of the expressions is in the \(\{\) path_expression : expr\} sequence is from left to right.
7. the expression returns the oid of the created object.

For instance:
- new Person() creates a person with all its attributes unitialized,
- Person() creates a person with all its attributes unitialized,
- new Person(name : "john") creates a person with its attribute name initialized to john,
- Person(name : "john") creates a person with its attribute name initialized to john,
- new Person(name : "john", age : 32, spouse : new Person(name : "mary")) creates a person named mary and a person named john, age 32 whose spouse is the person mary.

The new operator can also be used to create basic type object. Note that in this case, the operator is mandatory. The syntax for basic type creation is as follows:
new \([<[\) expr \(]>]\) basic_type (value).
1. where basic_type denotes an ODL basic type. It may be one of the following type: int32, int16, int32, char, byte, float or oid. Note that the type string is not allowed here.

\section*{5. LANGUAGE SYNTAX}
2. the value must be an atomic value of an OQL type mapped from the ODL basic type

For instance:
- new int(2)
- new float(2.3)
- new float(2)
- new char('a')
- new oid(first(select Person))

Finally, the new operator can be also used to create collections. The syntax for collection creation is as follows: [new] [<[expr]>] coll_type< class_name [, coll_name]> ([collection of elements])
1. the new operator is optionnal,
2. where coll_type denotes type of the collection: set, bag, array or list,
3. the class_name denotes the name of the class of the elements of the collection, for instance Person*, Car*,
4. coll_name is an optionnal string which denotes the name of the collection to create,
5. the optionnal collection of elements within parenthesis contains the elements (generally oids) to insert initially in the created collection,
6. the expression returns the oid of the created collection.

For instance:
- new set<Person *>() creates an empty set of persons,
- new set<Person *>(list(select Person)) creates a set containing all the persons in the database,
- new set<Person *, "all babies">(list(select Person.age < 1)) creates a set named all babies containing all the persons whose age is less than 1.
- new array<Car *>() creates an empty array of cars.
- new array<int>(list(1, 2, 3, 4) creates an array of integers initially containing \(1,2,3,4\).
- new set<int *> (list(new int(2), new int(10))) creates a set of integer objects containing initially two integer objects.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & new \\
\hline Syntax & new \([<[\) expr \(]>]\) class_name( \(\{\) path_expression : expr \(\}\) ) \\
\hline Type & n-ary \\
\hline Operand Types & any type \\
\hline Result Type & oid or object \\
\hline Functions & creates an persistent or transient object \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|l|}{Expression Examples} \\
\hline expression & result \\
\hline \[
\begin{aligned}
& \text { john }:=\text { new Person(name: "john", } \\
& \text { lastname: "wayne", } \\
& \text { age : fib(10)); }
\end{aligned}
\] & returns the oid of the created Person instance \\
\hline
\end{tabular}
\begin{tabular}{|c|l|}
\hline & \begin{tabular}{l} 
returns the oid of the created \\
new Person(name: "mary", \\
lastname: "poppins", \\
addr.town : "jungle", \\
addr.street[0] : 'a', \\
addr.street[1] : 'b', \\
spouse : john, \\
spouse.age \(: 72\) \\
);
\end{tabular} \\
\hline
\end{tabular}

\subsection*{5.23 Object Deletion}

The delete unary operator is used to delete persistent objects.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & delete \\
\hline Syntax & delete expr \\
\hline Type & unary \\
\hline Operand Types & oid or object \\
\hline Result Type & the operand type \\
\hline Functions & delete a transient or persistent object \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline expression & result \\
\hline delete first(select Person) & \begin{tabular}{l} 
the oid of the deleted Person in- \\
stance
\end{tabular} \\
\hline delete new Person() & \begin{tabular}{l} 
the oid of the deleted Person in- \\
stance
\end{tabular} \\
\hline for ( x in (select Person) delete x & \begin{tabular}{l} 
the oids of the deleted Person in- \\
stances
\end{tabular} \\
\hline
\end{tabular}

\subsection*{5.24 Collection Expressions}

OQL introduces a few operators for object collection manipulation: one of them is the array deferencing operator "[]" (5.15) that is overloaded for ordered collection manipulation.

Some them are ODMG OQL compliant, the others are EyEDB extensions. Object collections may be persistent or transient, orderered or not. These operators allows us to make the following kind of operations:
1. gets the contents of a collection: operator contents,
2. get an element at a given position in an ordered collection: operator [] (ODMG compliant),
3. get elements at some given positions in an ordered collection: operators [:] (ODMG compliant) and [?],
4. checks if an element is in a collection: operator in (ODMG compliant),
5. add an element in an unordered collection: operator add/to
6. suppress an element from an unordered collection: operator suppress/from,
7. set or suppress an element in an ordered at a given position: operator [],
8. set or suppress elements in an ordered at a given position: operators [:] and [?],
9. append an element in an ordered collection: operator append/to,
10. checks if a given condition is realized for at least one element in a collection: operator in (ODMG compliant),
11. checks if a given condition is realized for all elements in a collection: operator for/all (ODMG compliant),
12. checks if a given condition is realized for a given number range of elements in a collection: operator extended for.

In all the following examples, the OQL variables p0 denotes the first Person instance in the database: p0 := first (select Person).

Important Note: although they are reference in the following descriptions of collection operators, the object collections list are not implemented in the current version of EyEDB.

\section*{contents Operator}

The contents unary operator is used to give the contents of a given ordered or unordered object collection. It returned an OQL collection of the same type of the object collection: set, bag, array or list.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & contents \\
\hline Syntax & contents expr \\
\hline Type & unary \\
\hline Operand Types & oid or object collection \\
\hline Result Type & a collection of objects \\
\hline Functions & returns the contents of an object collection \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline contents (p0.children) & an array of Person oids \\
\hline select contents(x.children) from Person x & \begin{tabular}{l} 
returns a list of arrays of Person \\
oids.
\end{tabular} \\
\hline contents (list(1, 2, 3)) & \begin{tabular}{l} 
raises an error: oid or object \\
expected, got list
\end{tabular} \\
\hline
\end{tabular}

\section*{in Operator}

The in operator is used to check if a given element is in ordered or unordered object collection.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & in \\
\hline Syntax & expr in expr \\
\hline Type & binary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: any type, second operand: oid or object \\
collection
\end{tabular} \\
\hline Result Type & boolean \\
\hline Functions & \begin{tabular}{l} 
returns true if the first operand belongs to the collection \\
pointed by the second operand; false otherwise
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline \begin{tabular}{l} 
first(select Person.name \(=\) NULL) in \\
p0.children
\end{tabular} & \begin{tabular}{l} 
returns true if the first Person in- \\
stance whose name is unitialized is \\
in the array of children of the first \\
Person
\end{tabular} \\
\hline \begin{tabular}{l} 
first(select Car.brand \(=\) "renault") in \\
p0.cars
\end{tabular} & \begin{tabular}{l} 
returns true if the first Car instance \\
whose brand equals renault is in \\
the set of cars of the first Person
\end{tabular} \\
\hline
\end{tabular}

\section*{add/to Operator}

The add/to operator is used to add an element in an unordered collection (set or bag).
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & add/to \\
\hline Syntaxes & add expr to expr \\
\hline Type & binary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: any type, second operand: oid or object \\
unorderered collection (set or bag)
\end{tabular} \\
\hline Result Type & type of the first operand \\
\hline Functions & \begin{tabular}{l} 
adds the first operand to the non-indexed collection (i.e. \\
bag or set) pointed by the second operand; returns the \\
first operand.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline add new Car(num : 100) to p0.cars & returns the created Car oid \\
\hline \begin{tabular}{l} 
add new Person(name : "john") to \\
p0.children
\end{tabular} & \begin{tabular}{l} 
raises an error: cannot used non \\
indexed insertion in an array
\end{tabular} \\
\hline add new Car() to new set<Car *>() & \begin{tabular}{l} 
returns the just created car; but we \\
have lost the oid of the just created \\
set of cars!
\end{tabular} \\
\hline \begin{tabular}{l} 
add new Person() to (c := new bag<Person \\
\(*>())\)
\end{tabular} & \begin{tabular}{l} 
returns the just created person; the \\
created bag has been kept in the \\
OQL variable c
\end{tabular} \\
\hline
\end{tabular}

\section*{[] Operator}

The polymorphic [] operator is used to set or get an element in an ordered collection (array or list) at a given position: it can be used in a right or left value.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & {[]} \\
\hline Syntaxes & expr \([\) expr \(]\) \\
\hline Type & binary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: collection array or list, second operand: \\
integer
\end{tabular} \\
\hline Result Type & the type of the element, \\
\hline Functions & \begin{tabular}{l} 
gets the element in the collection pointed by the first \\
operand at the position pointed by the second operarand. \\
If used at a left value, gets a reference to that element.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline p0.children[0] & \begin{tabular}{l} 
returns the child at position \#0 in \\
p0.children collection. Returns \\
nil if there is no child at this posi- \\
tion
\end{tabular} \\
\hline p0.children[0] := Person(name : "john") & returns the created Person oid \\
\hline \begin{tabular}{l} 
p0.children[12039] \(:=\) Person(name : \\
"henry")
\end{tabular} & \begin{tabular}{l} 
returns the created Person oid. \\
This expression is valid in any cas \\
as the collection arrays automati- \\
cally increased its size
\end{tabular} \\
\hline \begin{tabular}{l} 
(array<Person \(*>())[0]:=\) new Car(num : \\
\(100)\)
\end{tabular} & \begin{tabular}{l} 
returns the just created person; but \\
the created array has been "lost" \\
as it is not tied to any instance and \\
as we did not bind it to any OQL \\
variable
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \begin{tabular}{l}
\((c:=\operatorname{array<Person} *>())[0]:=\) new \(\operatorname{Car}(\) num \\
\(: ~ 100)\)
\end{tabular} & \begin{tabular}{l} 
returns the just created person; the \\
created array has been kept in the \\
OQL variable c
\end{tabular} \\
\hline p0.cars [1] \(:=\operatorname{Car}(\) num : 100) & \begin{tabular}{l} 
raises an error: array expected, \\
got set
\end{tabular} \\
\hline
\end{tabular}

\section*{[:] Operator}

The polymorphic [:] operator is used to set or get some elements in an ordered collection (array or list) at some given positions: it can be used in a right or left value. When used in a right value, the returned atom is a set of struct with the two attributes index and value. In each struct element returned, the value of index is the position of the element, the value of value is the element value. Note the returned struct elements are not ordered according to the element postions; it is why a set is returned. When used as a left value, the returned atom is a set of references on the elements .
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & {\([:]\)} \\
\hline Syntaxes & \(\operatorname{expr}[\operatorname{expr}: \operatorname{expr}]\) \\
\hline Type & ternary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: collection array or list, second and third \\
operands: integer
\end{tabular} \\
\hline Result Type & the type of the element, \\
\hline Functions & \begin{tabular}{l} 
gets the elements in the collection pointed by the first \\
operand at the position range pointed by the second and \\
third operarands. If used at a left value, gets references \\
to that elements.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|l|}{Expression Examples} \\
\hline \[
\text { p0.children }[0: 1]
\] & returns a set of struct including the children and the position of the children position \(\# 0\) and \(\# 1\) in the p0.children collection. For instance: set(struct(index : 0, value : 3874.33.293847:oid), struct(index : 1, value : 2938.33.1928394:oid)) \\
\hline \multicolumn{2}{|l|}{Returns nil if there is no child at these positions} \\
\hline \[
\begin{aligned}
& \text { p0.children }[0: 4]:=\text { Person(name : "john") } \\
& \text { returns the created Person oid }
\end{aligned}
\] & Sets all the children at the position \(\# 0\) to \(\# 4\) to a new Person instance. \\
\hline ```
p0.children[12000:12039] := Person(name :
"henry")
``` & returns the created Person oid. This expression is valid in any cas as the collection arrays automatically increased its size \\
\hline (array<Person *>(list(Person())) [0] & returns the just created person within the just created array. But the array is "lost" as it is not tied to any instance and as we did not bind it to any OQL variable \\
\hline (x := array<Person *>(list(Person()))) [0] & returns the just created person; the created array has been kept in the OQL variable c \\
\hline
\end{tabular}

\section*{[?] Operator}

The polymorphic [?] operator is used to set or get all the elements in an ordered collection (array or list). It can be used in a right or left value. When used in a right value, the returned atom is a set of struct with the two attributes index and value. In each struct element returned, the value of index is the position of the element, the value of value is the element value. Note the returned struct elements are not ordered according to the element postions; it is why a
set is returned.
When used as a left value, the returned atom is a set of references on the elements .
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & {\([?]\)} \\
\hline Syntaxes & expr [?] \\
\hline Type & unary \\
\hline Operand Type & collection array or list, \\
\hline Result Type & a set of struct or a set of references \\
\hline Functions & \begin{tabular}{l} 
gets all the elements in the collection pointed by the first \\
operand If used at a left value, gets references to that \\
elements.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|l|}{Expression Examples} \\
\hline p0.children [?] & returns a set of struct including the children and the position of all the children in the po.children collection. For instance: set(struct(index : 0, value : 3874.33.293847:oid), struct(index : 1, value : 2938.33.1928394:oid)) \\
\hline \multicolumn{2}{|l|}{Returns nil if the collection is empty} \\
\hline ```
p0.children[?] := Person(name : "john")
returns the created Person oid
``` & Sets all the children to a new Person instance. \\
\hline ```
(array<Person *>(list(Person(),
Person()))) [?]
``` & returns a set of struct including the just created Person instances in the just created array. \\
\hline
\end{tabular}

\section*{append/to Operator}

The append/to operator is used to append an element to an ordered collection (list or array).
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & append \\
\hline Syntaxes & append expr to expr \\
\hline Type & binary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: any type, second operand: oid or object \\
denoting an ordered collection
\end{tabular} \\
\hline Result Type & any type \\
\hline Functions & \begin{tabular}{l} 
appends the element denoted by the first operand to the \\
indexed collection (i.e. list or array) denoted by the sec- \\
ond operand.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline append Person() to p0.children & the created Person instance \\
\hline append Car()to p0.cars & \begin{tabular}{l} 
raises an error: array or list \\
expected, got set<Person*>
\end{tabular} \\
\hline
\end{tabular}

\section*{suppress/from Operator}

The suppress/from operator is used to suppress an element from an ordered collection (set or bag).
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & suppress/from \\
\hline Syntaxes & suppress expr from expr \\
\hline Type & binary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: any type, second operand: oid or object \\
collection
\end{tabular} \\
\hline Result Type & type of the first operand \\
\hline Functions & \begin{tabular}{l} 
suppress the first operand from the non-indexed collec- \\
tion (i.e. bag or set) pointed by the second operand; \\
returns the first operand.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline \begin{tabular}{l} 
suppress (select Car.num \(=1000\) ) from \\
p0.cars
\end{tabular} & \begin{tabular}{l} 
the suppressed car if it was found in \\
the collection; otherwise, raises an \\
error
\end{tabular} \\
\hline suppress new Car() from p.cars & \begin{tabular}{l} 
raises an error: item \\
'71238.13.3959935:oid' not \\
found in collection
\end{tabular} \\
\hline suppress p0 from p0.children & \begin{tabular}{l} 
raises an error: cannot used \\
non indexed suppression in an \\
array
\end{tabular} \\
\hline
\end{tabular}

\section*{empty Operator}

The empty operator is used to empty an ordered or an unordered collection.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & empty \\
\hline Syntax & empty expr \\
\hline Type & unary \\
\hline Operand Types & oid or object collection \\
\hline Result Type & nil \\
\hline Functions & empty the collection pointed by the operand \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline empty(first (select Person).children) & nil \\
\hline empty(first (select Person).cars) & nil \\
\hline empty new set<Car *>(list(new Car())) & \begin{tabular}{l} 
nil; this expression creates a col- \\
\\
\\
\\
\\
\\
\\
\end{tabular} nection of Car containing initially a and empty it! \\
\hline
\end{tabular}

\section*{in Operator}

The in operator is used to check if a given condition is realized for at least one element in an ordered or unordered collection.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & in \\
\hline Syntax & identifier in expr : expr \\
\hline Type & ternary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: identifier, second operand: oid or \\
object collection, third operand: boolean
\end{tabular} \\
\hline Result Type & boolean \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline Functions & \begin{tabular}{l} 
returns true if it exists in the collection pointed by the \\
second operand an element for which the third operand \\
is evaluated to true.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|l|}
\hline \multicolumn{3}{c|}{ Expression Examples } \\
\hline x in p0.children: x. name \(=\) "mary" & true or false \\
\hline x in p0.cars: x. num \(<100\) and x. num \(>=90\) & true or false \\
\hline
\end{tabular}

\section*{for Operator}

The for/all operator is used to check if a given condition is realized for all elements in an ordered or unordered collection. This operator is ODMG compliant.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & for/all \\
\hline Syntaxes & for all identifier in expr : expr \\
\hline Type & ternary \\
\hline Operand Types & \begin{tabular}{l} 
first operand: identifier, second operand: oid or \\
object collection, third operand: boolean
\end{tabular} \\
\hline Result Type & boolean \\
\hline Functions & \begin{tabular}{l} 
returns true if for all items contained in the collection \\
pointed by the second operand the third operand is eval- \\
uated to true.
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline \begin{tabular}{l} 
for all x in p0.children: x. name \(==\) \\
"john"
\end{tabular} & true or false \\
\hline for all x in p0.cars: x.num \% 10 & true or false \\
\hline
\end{tabular}

The for/cardinality operator is used to check if a given condition is realized for a given number range of elements in an orderer or unordered collection. Note that this operator is the generalisation of the in and for/all operators:
for \(\langle 0: \$\rangle\) is equivalent to in
for \(\langle \$>\) is equivalent to for all
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & for cardinality \\
\hline Syntaxes & \begin{tabular}{l} 
for < expr : expr > identifier in expr : expr \\
for < expr > identifier in expr : expr
\end{tabular} \\
\hline Type & 5 -ary \\
\hline Operand Types & \begin{tabular}{l} 
first and optional second operands: integer or \$, where \\
\$ denotes the collection cardinality, third operand: oid \\
or object, fourth operand: identifier, fifth operand: \\
boolean
\end{tabular} \\
\hline Result Type & boolean \\
\hline Functions & \begin{tabular}{l} 
returns true if the number of items in the collection \\
pointed by the third operand for which the third operand \\
is evaluated to true is in the interval [first operand, sec- \\
ond operand].
\end{tabular} \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline \begin{tabular}{l} 
for \(\langle 0: 4\rangle \mathrm{x}\) in p0.children: x.name \(==\) \\
"john"
\end{tabular} & \begin{tabular}{l} 
true if at most 4 children have their \\
name equals to john
\end{tabular} \\
\hline \begin{tabular}{l} 
for \(\langle 4\rangle \mathrm{x}\) in p0.children: x. name \(==\) \\
"john"
\end{tabular} & \begin{tabular}{l} 
true if one an only one children \\
have its name equals to john
\end{tabular} \\
\hline for <0:\$> x in p0.cars: x.num \(=10\) & equivalent to in \\
\hline for \(\langle \$\rangle \mathrm{x}\) in p0.cars: x. num \(=10\) & equivalent to for/all \\
\hline
\end{tabular}

\subsection*{5.25 Exception Expressions}

Currently, EyEDB OQL does not provide full support for exception management as there is no try/catch operator. Nevertheless, the throw operator allows us to raise an error message, for instance:
```

if (!check(p))
throw "variable p is not correct".

```

The throw operator stops the current thread of statements and returns the error message at the uppest level. In the following code:
```

a := 1;
throw "this is an error";
b := 2;

```
the variable a will be assigned to 1 , but the variable b will not be assigned to 2 as the throw expression deroutes the normal thread of statements. The throw operator is often used in the body of functions,
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & throw \\
\hline Syntax & throw expr \\
\hline Type & unary \\
\hline Operand Type & string \\
\hline Result Type & nil \\
\hline Functions & raises an error message \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline throw "error message" & nil \\
\hline throw "error \#1: " +msg & nil \\
\hline
\end{tabular}

\subsection*{5.26 Function Definition Expressions}

As introduced previously, OQL supports functions. There are two types of functions definition syntax: function definition expression and function definition statements. The first ones, exposed in this section, are more restrictive than the second ones, as their definition can contain only one expression. The second ones contain an unbounded sequence of statements. The function definition expressions are ODMG compatible and are called Named Query Definition in the ODMG standard. To define such a function, one must use the operator define/as. The general syntax for a definition function expression is:
define identifier \([(\) arglist \()]\) as expr.
For instance:
define \(\operatorname{div} 2(\mathrm{x})\) as \(\mathrm{x} / 2\);
div2(10); // returns 5
define pimul(x) as \(\mathrm{x} * 3.1415926535\);
pimul(23); // returns 72.256631
define getOnePerson(name) as first(select Person.name = name);
getOnePerson("john"); // returns an oid or nil

As the last operand of the define/as operator is any OQL expression, it can be a sequence of expressions by using the comma sequencing operator. Therefore, the following construct is valid:
```

define getit(age) as ::getit_call++,
list_of_persons := select Person.age >= age,
(count(list_of_persons) > 0 ? list_of_persons[0] : nil);
getit(10); // returns the first Person whose age is greater or equal to
10, or nil
getit_call; // equals 1
list_of_persons; // raises an error: uninitialized identifier 'list_of_persons
getit(20);
getit_call; // equals 2

```

Several comments about this code:
1. : :getit_call denotes a global variable, while list_of_persons denotes a variable local to the function: the variable scoping has previously been introduced in the identifier expressions, and will be explained again in the function definition statements Section.
2. as there are no iteration expression (for instance a for or while expression) and as a function definition expression can contain only one expression, one cannot use a function definition expression with iteration statements. To use an iteration statement, one needs to use a function definition statement.
3. when one needs to define a function with a sequence of expressions, it may be easier and more readable to use a function definition statement instead of a function definition expression. The statement version of the previous function is:
```

function getit(age) {
::getit_call++;
list_of_persons := select Person.age >= age;
if (count(list_of_persons))
return list_of_persons[0];
}

```
which is - for C, C++ or Java programmers - a more natural construct.
Finally, the function definition allows us for recursive definitions, for instance:
```

define fib(n) as (n < 2 ? n : fib(n-2) + fib(n-1));
fib(10); // returns 55

```
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & define/as \\
\hline Syntax & define identifier [arglist] as expr \\
\hline Type & n-ary \\
\hline Operand Type & \begin{tabular}{l} 
first operand: identifier, last operand: any type, other \\
operands: identifier
\end{tabular} \\
\hline Result Type & identifier \\
\hline Functions & defines a function \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{3}{|c|}{ Expression Examples } \\
\hline define Persons as select Person & Persons \\
\hline \begin{tabular}{l} 
define fact \((\mathrm{n})\) \\
fact \((\mathrm{n}-1))\)
\end{tabular} & fact \\
\hline
\end{tabular}

\subsection*{5.27 Conversion Expressions}

The conversion unary operators string, int, char, float, oid and ident allows us to make the following conversions:
\begin{tabular}{|l|l|l|l|}
\hline Operator & From & To & Returned Atom \\
\hline string & any type & string & the string representation of the operand \\
\hline int & \begin{tabular}{l} 
int \\
char \\
float \\
string
\end{tabular} & \begin{tabular}{l} 
int \\
int \\
int \\
int
\end{tabular} & \begin{tabular}{l} 
the int operand \\
the operand casted to an int \\
the operand casted to an int \\
the operand converted to an int
\end{tabular} \\
\hline char & \begin{tabular}{l} 
char \\
int \\
float \\
string
\end{tabular} & \begin{tabular}{l} 
char \\
char \\
textttchar \\
char
\end{tabular} & \begin{tabular}{l} 
the char operand \\
the operand casted to a char \\
the operand casted to a char \\
the operand converted to a char
\end{tabular} \\
\hline float & \begin{tabular}{l} 
float \\
char \\
int \\
string
\end{tabular} & \begin{tabular}{l} 
float \\
float \\
float \\
float
\end{tabular} & \begin{tabular}{l} 
the float operand \\
the operand casted to a float \\
the operand casted to a float \\
the operand converted to a float
\end{tabular} \\
\hline oid & \begin{tabular}{l} 
oid \\
string
\end{tabular} & \begin{tabular}{l} 
oid \\
oid
\end{tabular} & \begin{tabular}{l} 
the oid operand \\
the string operand converted to an oid
\end{tabular} \\
\hline ident & \begin{tabular}{l} 
ident \\
string
\end{tabular} & \begin{tabular}{l} 
ident \\
ident
\end{tabular} & \begin{tabular}{l} 
the ident operand \\
the string operand converted to an ident
\end{tabular} \\
\hline
\end{tabular}

These operators are used to perform an explicit conversion such as convert the string "123" to the integer 123, or to perform an explicit cast for numbers such as casting the integer 10 to the float 10.0. These operators evaluate first their operand before performing the conversion. If the operand type is valid, no error is raised even if its format is not valid, for instance: int "alpha" returns 0, while oid "aoaoai" returns NULL. Note that because of the precedence of these operators, parenthesis are necessary to make a conversion of a non-primary operand. For instance, string \(1+2\) is not valid: you should use string (1+2).

\section*{string operator}

The string operator evaluates its operand and returns its string representation.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & string \\
\hline Syntax & string expr \\
\hline Type & unary \\
\hline Operand Type & any type \\
\hline Result Type & string \\
\hline Function & returns the string representation of any atom \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline string 123.3 & "1203.300000" \\
\hline string 'a' & "a" \\
\hline string first(select Person) & "71211.13.1486847: oid" \\
\hline string \&alpha & ": :alpha" \\
\hline string list(1, 2, 3+2) & "list(1, 2, 5)" \\
\hline string (list("hello", 30) + list(10)) & "list("hello", 30, 10)" \\
\hline string (1+3) & "4" \\
\hline string 1+3 & raises an error \\
\hline
\end{tabular}

\section*{int operator}

The int operator evaluates its operand and converts or casts it to an integer.
If the operand is the string, it converts it using the atoi \(C\) function. If the string is not a valid integer, it returns a 0 . If the operand is a char or float, it casts it to an integer.
If the operand is an integer, it returns it.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & int \\
\hline Syntax & int expr \\
\hline Type & unary \\
\hline Operand Type & int, char, float or string \\
\hline Result Type & int \\
\hline Function & returns the integer conversion or cast of the operand \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline int 123.3 & 123 \\
\hline int 12 & 12 \\
\hline int 'a' & 97 \\
\hline int "123" & 123 \\
\hline int ("123" + "12") & 12312 \\
\hline int alpha & \begin{tabular}{l} 
the value of alpha converted or \\
casted to an integer
\end{tabular} \\
\hline int list \((1,2,3)\) & raises an error \\
\hline
\end{tabular}

\section*{char operator}

The char operator evaluates its operand and converts or casts it to a char.
If the operand is the string of length one, it returns the character of this string. If the string has several characters, it returns a ' \(\backslash 000\).
If the operand is a integer or float, it casts it to a character.
If the operand is a character integer, it returns it.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & char \\
\hline Syntax & char expr \\
\hline Type & unary \\
\hline Operand Type & int, char, float or string \\
\hline Result Type & char \\
\hline Function & returns the character conversion or cast of the operand \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline char 123.3 & \(\{\) \\
\hline char 'a' & 'a' \\
\hline char alpha & \begin{tabular}{l} 
the value of alpha converted or \\
casted to a character
\end{tabular} \\
\hline char "a" & 'a' \\
\hline char "hello" & '^@' \\
\hline char list(1, 2, 3) & raises an error \\
\hline
\end{tabular}

\section*{float operator}

The float operator evaluates its operand and converts or casts it to a float.
If the operand is the string, it converts it using the atof \(C\) function. If the string is not a valid float, it returns a 0.0 .
If the operand is a integer or float, it casts it to a float.
If the operand is a float, it returns it.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & float \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline Syntax & float \(\operatorname{expr}\) \\
\hline Type & unary \\
\hline Operand Type & int, char, float or string \\
\hline Result Type & float \\
\hline Function & returns the float conversion or cast of the operand \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline float 123.0 & 123.0 \\
\hline float 123.3 & 123.3 \\
\hline float 'a' & 97.000 \\
\hline float "123.0000000" & 123.0 \\
\hline float ("123." + "12") & 123.12 \\
\hline float "hello" & 0.0 \\
\hline float alpha & \begin{tabular}{l} 
the value of alpha converted or \\
casted to a float
\end{tabular} \\
\hline float list(1, 2, 3) & raises an error \\
\hline
\end{tabular}

\section*{oid operator}

The oid operator evaluates its string operand and returns the corresponding oid. If the string does not denote a valid oid, the NULL oid is returned.
If the operand is an oid, it returns it.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & oid \\
\hline Syntax & oid expr \\
\hline Type & unary \\
\hline Operand Type & oid or string \\
\hline Result Type & oid \\
\hline Function & returns the oid denoted by the string operand \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline oid "234.34.33:oid" & \(234.34 .33:\) oid \\
\hline oid 234.34.33:oid & \(234.34 .33:\) oid \\
\hline oid first (select Person) & returns the first person oid \\
\hline oid 'a' & raises an error \\
\hline
\end{tabular}

\section*{ident operator}

The ident operator evaluates its string operand and returns the corresponding identifier. If the operand is an identifier, it returns it.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & ident \\
\hline Syntax & ident expr \\
\hline Type & unary \\
\hline Operand Type & ident or string \\
\hline Result Type & string \\
\hline Function & returns the identifier denoted by the string operand \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline ident "alpha" & alpha \\
\hline ident "alpha\#1x" & alpha\#1x \\
\hline ident "alpha" := 123 & 123, alpha has been assigned to 123 \\
\hline valof \&(ident "alpha") & 123 \\
\hline ident 'a' & raises an error \\
\hline
\end{tabular}

\subsection*{5.28 Type Information Expressions}

OQL provides two type information unary operators: typeof and classof. The first one takes any operand type, while the second one takes an oid or an object operand. Note that because of the precedence of these operators, parenthesis are necessary to get type information about a non-primary operand. For instance, typeof \(1+2\) is not valid: you should use typeof (1+2).

\section*{typeof operator}

The typeof operator is used to get the type of any OQL atom. It evaluates its operand and returns the string type of its operand. For instance: typeof 1 returns "int" while typeof "hello" returns "string".
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & typeof \\
\hline Syntax & typeof expr \\
\hline Type & unary \\
\hline Operand Type & any type \\
\hline Result Type & string \\
\hline Function & returns the type of the atom \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline typeof "alpha" & "string" \\
\hline typeof (1+20.) & "float" \\
\hline typeof list(1, 2, 3) & "list" \\
\hline typeof first(select Person) & "oid" \\
\hline typeof 1+3049 & raises an error \\
\hline typeof alpha & type of the value of alpha \\
\hline typeof \&alpha & ident \\
\hline
\end{tabular}

\section*{classof operator}
typeof operator
The classof operator is used to get the class of any oid or object. It evaluates its operand and returns the string class of its operand. For instance: classof first (select Person) returns "Person" while typeof new Car() returns "Car".
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & classof \\
\hline Syntax & classof expr \\
\hline Type & unary \\
\hline Operand Type & oid or object \\
\hline Result Type & string \\
\hline Function & returns the class of the operand \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline classof first(select class) & "basic_class" \\
\hline classof (c := new Car (num : 10)) & "Car" \\
\hline classof NULL & "" \\
\hline classof first(select Person).spouse & "Person" or NULL \\
\hline classof 1 & raises an error \\
\hline
\end{tabular}

\subsection*{5.29 Query Expressions}

The select operator is used to perform queries in a database. The general syntax of this operator is as follows:
select [distinct] projection [from fromList [where predicat] [order by orderExprList [asc|desc]]
1. distinct means that duplicates must be eliminated,
2. projection is an expression using the variables defined in the fromList,
3. fromList is a sequence of comma-separated items under one of the following forms:
var in expr
expr as var
expr var
where expr is an expression of type collection or is a class name, and var is the name of a variable.
4. predicat is a boolean expression using the variables defined in the fromList,
5. orderExprList is a comma-separated list of sortable expressions (i.e. atomic type).
6. asc means that the order should be performed in an ascendant way (the default) and desc means the inverse.

\section*{ODMG vs. EyeDB OQL Query Expressions}

As explained in the Section \(O Q L\) vs. \(O D M G 3 O Q L\), there are a few differences between ODMG OQL and EyEDB OQL query expressions:
- the having/group clause is not supported in the current implementation.
- in the current implementation, one cannot use implicit from clause (i.e. from clause without variables). ODMG OQL supports constructs such as: from Person without any variable. This implementation does not.
- contrary to ODMG OQL, the from clause is optionnal, A select expression which does not use the from is called an implicit select expression.
- in a from clause such as "x in expr", "expr" can be the name of a class. In this case, the interpreter understand this as the extent of this class. ODMG OQL does not support that.
- the SQL specific agregate operators \(\min (*), \max (*)\), count \((*), \operatorname{sum}(*)\) and \(\operatorname{avg}(*)\) are not supported
- the order by clause is more restrictive than in the ODMG OQL specifications (see below).
- the select \(*\) is not supported in the current implementation.

\section*{The general select syntax}

Rather than introducing the select operator in a formal way by using a lot mathematical symbols, we introduce it first, in an unformal way, and then, through query examples.

The unformal description of the query process is as follows:
1. The from clause determine the sets of objects on which the query will be applied. For instance, in the from clause, "x in Person, y in \(x\).children", the sets on which the query will be applied are all the Person instances bound to the variable \(x\) and the children of these instances, bound to the variable \(y\).
2. These sets of objects are filtered by retaining only the objects that satisfy the predicat in the where clause. These result objects are gathered into a bag. If no where clause is there, all objects are retained.
3. If an order by clause is present, a sort is performed using to the following process:
(a) each order expression must be of a sortable type: number (int, char or float) or string. If not, an error is raised.
(b) the bag is ordered into a list according to the first order expression. Then identical atoms are ordered again using the second order expression and so on.
(c) there is a restriction in the current implementation: each expression in the orderExprList must be present in the projection expression. If not present, an error is raised.
4. The projection expression is evaluated for each object in the collection and the results of these evaluations are gathered into a bag.
5. If the keyword distinct is there, the eventual duplicates are eliminated.
6. Finally, if the order by clause is present, the result bag is converted to a list; if the distinct keyword is there without an order by, the bag is converted to a set; and if neither order by nor distinct are used, we get a bag.

The following table presents several examples:
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|r|}{Simple select/from Examples} \\
\hline \multicolumn{2}{|r|}{\begin{tabular}{l}
select x from Person x \\
returns a bag containing all the Person oids in the database
\end{tabular}} \\
\hline \multicolumn{2}{|l|}{select x.name from Person \(x\) returns a bag containing the name of every Person instance in the database} \\
\hline \multicolumn{2}{|l|}{select struct (name: x.name, age: x.age) from Person \(x\) returns a bag containing struct elements including the name and age of every Person instance in the database} \\
\hline \multicolumn{2}{|l|}{select list(x, x.name) from Person \(x\) where x.name ~ "h" returns a bag of list elements containing the oid and the name of every Person instances whose name matches the regular expression "h"} \\
\hline \multicolumn{2}{|l|}{select list( \(x, x . n a m e)\) from Person \(x\) where \(x . n a m e\) ~ "h" order by x.name same as previous example, but the result is a list ordered by the name of the persons} \\
\hline \multicolumn{2}{|l|}{select \(x\) from Person \(x\) where \(x . s p o u s e . n a m e=" j o h n "\) or \(x . a g e<10\) returns a bag of Person instances whose spouse name is equal to "john" or the age is less than 10} \\
\hline \multicolumn{2}{|l|}{```
select x from Person x order by x.name
current implementation restriction: raises an error: x.name not found in
projection
```} \\
\hline
\end{tabular}

\section*{Arrays and collections in query expressions}

OQL provide supports for performing direct queries through non-collection array and collection attributes without explicit joins. To perform such queries, one must use the operators [], [?] or [:]. All the operators may be used for non-collection arrays. For collections (list, set, bag and array), only the operator [?] is valid.
The operator [] denotes an element in a non-collection array.
The operator [:] denotes a range of elements in a non-collection array.
The operator [?] denotes all elements in a non-collection array or in a collection. The following table presents several examples:

\begin{tabular}{|l|} 
returns all the persons whose one of its children is called "johnny" \\
\hline select \(x\) from Person \(x\) where \(x . c a r s[?] . n u m ~<~\) \\
"mary" \\
returns all the persons whose one of its cars has a number less than 100 or a child \\
called "mary" \\
\hline select x from Person x where x. children[1]. name \(=\) "johnny" \\
although the children is a collection array, an error is raised. This is a current \\
limitation of the implementation that will disapear soon
\end{tabular}

\section*{The Implicit select syntax}

An implicit select expression is a select expression with neither a from nor an explicit where clause. In fact, the where clause may be included in the projection expression.
This particular syntax has the advantage to be more compact and more simple than the general syntax, but some queries cannot be performed:
1. queries performing an explicit join,
2. queries having or or and in their where clause.

The following table presents several examples:
\begin{tabular}{|l|}
\hline \\
\\
\hline \begin{tabular}{l} 
select 1 \\
returns 1
\end{tabular} \\
\hline \begin{tabular}{l} 
select Person \\
returns a bag containing all Person instances in the database
\end{tabular} \\
\hline \begin{tabular}{l} 
select Person.name \\
returns a bag containing the name of every Person instances in the database
\end{tabular} \\
\hline \begin{tabular}{l} 
select Person.name \(=\) "john" \\
returns a bag containing the oids of every Person instances whose name is equal to \\
"john"
\end{tabular} \\
\hline \begin{tabular}{l} 
(select distinct Person.name \(=\) "john"). age \\
returns a set containing the age of every Person instances whose name is equal to \\
"john"
\end{tabular} \\
\hline \begin{tabular}{l} 
select Person.name \(=\) "john" or Person. age \(=10\) \\
raises an error: use select/from/where clause
\end{tabular} \\
\hline \begin{tabular}{l} 
select Person.name order by Person.name \\
returns a list containing of the sorted names of all Person instances
\end{tabular} \\
\hline
\end{tabular}

\section*{Querying the schema}

As every abstraction is an object, one can perform queries on the schema and classes. For instance, to get the oids of all the existing classes in a schema:
select schema which is equivalent to select class.
In the EyEDB object model, the class class has the following native attributes (some are inherited from the object class):
1. class (inherited from object) is the class of this class,
2. protection (inherited from object) is the protection object of this class,
3. type is the type of the class: "system" or "user".
4. name is the name of this class,
5. parent is the parent class of this class,
6. extent is the extent collection of this class: contains all the object instances of this class,
7. components is the collection of components of this class: contains the constraints, the methods, the triggers, the index.

Queries can be performed according to one or more of the previous attributes.
\begin{tabular}{|l|}
\hline \\
\(\qquad\)\begin{tabular}{l} 
Query Schema Examples
\end{tabular} \\
\hline \begin{tabular}{l} 
select class.name \(=\) "Person" \\
returns a bag containing the class whose name is "Person"
\end{tabular} \\
\hline \begin{tabular}{l} 
select class.type \(=\) "user" \\
returns all the user classes
\end{tabular} \\
\hline \begin{tabular}{l} 
select x from class x where x.name \(\sim\) \\
returns the user classes whose name matches the given regular expression
\end{tabular} \\
\hline \begin{tabular}{l} 
select class.parent.name \(=\) "Person" \\
returns a bag containing the sub-classes of the class Person
\end{tabular} \\
\hline
\end{tabular}

\section*{How queries are optimized?}

EyEDB queries implementation make an heavy use of index. Index may be used as terminal index or as non-terminal index: terminal index are those used at the end of a path expression, and the other are those used inside a path expression.

For instance, assuming that each attribute in our schema is indexed, the query select Person.name = "john" uses the index on the attribute name as a terminal index.

The query select Person. spouse. age < 23 uses the index on the attribute age as a terminal index, and the index on the indirect attribute spouse as a non-terminal index, while the query select Person. spouse \(=6252.3 .48474\) :oid uses the index on the spouse attribute as a terminal index.

The query select Person.children [?].spouse.name = "mary" uses the index of the attributes name, spouse and of the literal collection attribute children.

The queries with a where clause containing logical and constructs are optimized so to take advantage of the index. For instance, assuming that there is an index on the name attribute and no index on the age attribute, the query select x from Person x where x .age \(=100\) and x .name \(=\) "john" will perform first the query on the name attribute and then filter the result according to the given predicat on the age attribute.
If the two expressions around the logical and operator have an index or if none of the two expression has a index, the interpreter performs first the left part and then the second part. So, in this case, the order of the two expressions around the and is important.

Finally, the query select Person. name reads directly the index of the name attribute, instead of reading the name of each Person instance.

Nevertheless, currently, there are some constructs that the OQL interpreter does not interpret cleverly and where index are not used. This is unfortanely the case of the join constructs such as:
select x from Person x , x .spouse y where y. name \(=\) "mary".
This construct will not make use of the index on the spouse attribute (but it will use the index of the name attribute). Fortunately, in a lot of cases, join queries may be replaced by a path expression query, for instance:
select \(x\) from Person \(x\) where \(x . s p o u s e . n a m e=\) "mary" is the alternate form the previous join query.
This alternate form (path-expression oriented) is the preferred one, because:
- it is more intuitive,
- it is more object oriented (please forget relationnal!),
- it is more compact,
- it uses index properly.

Of course, the last reason is not a good reason as a proper query implementation should use index whatever the used syntax. The implementation will be improved in a next version.

Note that query optimisations are fragile and do not believe that the OQL interpreter knows your data better than you. So, the two following simple rules should be applied:
1. when path expression oriented queries are enough, do not use join constructs,
2. in a and expression within a where clause, put the expression which denotes the most little set of instances on the left side of the and. In some particular cases, the OQL interpreter knows how to optimize the and, but in most of the cases, it does not.

Some experimental hints can be added to an and expression in the where clause to help the interpreter to do the things better, but there are currentlty too much experimental to be documented here.

\subsection*{5.30 Miscellenaous Expressions}

A few OQL operators cannot be easily classified in one of the previous categories. We choose to classify them in the miscellenaous operators. These operators are: bodyof, structof, [!] and import.

\section*{bodyof operator}

The bodyof operaror is used to get the body of an OQL function. For instance, let the function fib: define fib(n) as (n \(<2\) ? \(n: f i b(n-2)+f i b(n-1))\). The expression bodyof fib will return: "fib(n) ( \(n<2\) ) n : \((\mathrm{fib}((\mathrm{n}-2))+\mathrm{fib}((\mathrm{n}-1))))\). This operator could be applied to expression-functions (i.e. functions defined with the define/as operator) or statementfunctions (i.e. functions defined with the function operator).
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & bodyof \\
\hline Syntax & classof \(\operatorname{expr}\) \\
\hline Type & unary \\
\hline Operand Type & ident denoting a function \\
\hline Result Type & string \\
\hline Function & returns the body of the function \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline bodyof is_int & "is_int(x) ((typeof x)=="integer")" \\
\hline bodyof first & "first(l) \{ if (((!is_list(l))\&\&(!is_array(l)))) \\
& \begin{tabular}{l} 
return l; start:=0; f:=nil; for (x in l) if \\
\(((s t a r t==0)) ~\{~ s t a r t:=1 ; ~ f:=x ; ~ b r e a k ; ~\} ; ~ ; ~ r e t u r n ~\) \\
\(f ;\} "\)
\end{tabular} \\
\hline bodyof 1 & raises an error \\
\hline
\end{tabular}

\section*{structof operator}

The structof is used to get the meta-type of a struct atom. The meta-type of a struct atom is the list of the attributes of this struct. For instance, the meta-type of struct (a : 1, b: "hello") is the list composed of the two attribute a and b . The following expression structof \(\operatorname{struct(a~:~1,~b~:~"hello")~returns~list("a",~"b").~}\)
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & structof \\
\hline Syntax & structof \(\operatorname{expr}\) \\
\hline Type & unary \\
\hline Operand Type & struct \\
\hline Result Type & a list of strings \\
\hline Function & returns the meta-type of the operand \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline structof struct(alpha : 1, beta : 2) & list("alpha", "beta") \\
\hline \begin{tabular}{l} 
structof first(select struct(x: \\
x.firstname) from Person x\()\)
\end{tabular} & list("x") \\
\hline structof 1 & raises an error \\
\hline
\end{tabular}

\section*{[!] operator}

The [!] is used to get the length (or size) of an ordered or unordered collection, a string or a struct. For instance, "hello" [!] returns 5, while list \((1,2,3)[!]\) returns 3 . Note that this operator is far more efficient than the strlen
library function.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & {\([!]\)} \\
\hline Syntax & expr \([!]\) \\
\hline Type & unary \\
\hline Operand Type & string, collection or struct \\
\hline Result Type & a int \\
\hline Function & returns length of the operand \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline (select Person) [!] & \begin{tabular}{l} 
the number of person in- \\
stances
\end{tabular} \\
\hline (struct(a: 1, b:2, c: "hello")) [!] & 3 \\
\hline ("hello"+"world") [!] & 10 \\
\hline "hello"+"world"[!] & raises an error \\
\hline
\end{tabular}

\section*{import operator}

The import operator is used to import an OQL file in the current OQL session. Its operand is a string which denote the absolute path (i.e. beginning with a /) or relative path (i.e. not beginning with a/) of the file to import. When the path is relative, the OQL interpreter will look in every directories pointed by the EYEDB configuration variable oqlpath. By default, oqlpath is equal to \(\%\) root \(\% / \mathrm{etc} /\) oql. If the file name has no .oql extension, the OQL interpreter will automatically adds one.
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ General Information } \\
\hline Operator & import \\
\hline Syntax & import expr \\
\hline Type & unary \\
\hline Operand Type & string \\
\hline Result Type & a string \\
\hline Function & import the file \\
\hline
\end{tabular}
\begin{tabular}{|l|l|}
\hline \multicolumn{2}{|c|}{ Expression Examples } \\
\hline import "stdlib" & "/usr/local/eyedb/etc/so/stdlib.oql" \\
\hline import "roudoudou" & raises an error: cannot find file 'roudoudou' \\
\hline
\end{tabular}

\subsection*{5.31 Selection Statements}

The selection statement is based on the if/else constructs.
Syntax: if ( cond_expr ) statement1 [else statement2] where cond_expr is a boolean expression, and statement1 and statement2 may be any statement: an expression statement, an iteration statement, a compound statement, a selection statement, a function definition statement, a jump statement or an empty statement.

Semantics: if the boolean expression cond_expr is evaluated to true, the statement statement1 is executed. Otherwise, if an else part is there, the statement statement2 is executed. The statements 1 and 2 may be any statement: an expression statement, iteration statement, compound statement, selection statement, function definition statement, jump statement or empty statement.
Note that a selection statement does not return any atom

The following table presents several examples of if/else statements:


\subsection*{5.32 Iteration Statements}

The iteration statements are based on the following constructs:
- while
- do/while
- for C, C++, Java form
- for collection form

\section*{while statement}

Syntax: while ( cond_expr ) statement
where cond_expr is a boolean expression, and statement any statement: an expression statement, an iteration statement, a compound statement, a selection statement, a function definition statement, a jump statement or an empty statement.

Semantics: The statement is executed while the boolean expression cond_expr is evaluated to true. Note that a while statement does not return any atom

The following table presents several examples of while statements:
\begin{tabular}{|l|}
\hline \\
\\
\hline \begin{tabular}{l} 
while (true) a++; \\
a is increment definitively!
\end{tabular} \\
\hline \begin{tabular}{l} 
while (n--) a++; \\
an error is raised: boolean expected, got integer
\end{tabular} \\
\hline \begin{tabular}{l} 
while (n-- > 0) a++; \\
this is better
\end{tabular} \\
\hline \begin{tabular}{l} 
while (n++ <= \(100 ~|\mid ~ s t o p) ~\{i f ~(!p e r f o r m(a++)) ~ b r e a k ; ~ c h e c k(a) ~ ; ~\} ~\)
\end{tabular} \\
note the usage of a compound statement and of the break
\end{tabular}

\section*{do/while statement}

Syntax: do statement while ( cond_expr )
where cond_expr is a boolean expression, and statement any statement: an expression statement, an iteration statement, a compound statement, a selection statement, a function definition statement, a jump statement or an empty statement.

Semantics: The statement is executed at least once. Then while the boolean expression cond_expr is evaluated to true, the statement is executed. Note that a do/while statement does not return any atom

The following table presents several examples of do/while statements:
\begin{tabular}{|c|}
\hline do/while Statement Examples \\
\hline do a++; while (true); a is increment definitively! \\
\hline do a=+; while (n--); an error is raised: boolean expected, got integer \\
\hline do a++; while (n-- > 0); this is better \\
\hline do \{if (!perform(a++)) break; check(a);\} while (n++ <= 100 || stop); note the usage of a compound statement and of the break \\
\hline ```
do {l := (select Person.name = name); name := get_name();} while (name
!= "john");
no comments
``` \\
\hline
\end{tabular}

\section*{C-for statement}

Syntax: for ([expr1] ; [cond_expr] ; [expr2] ) statement
where cond_expr is a boolean expression, expr1 and expr2 are any expressionss and statement any statement: an expression statement, an iteration statement, a compound statement, a selection statement, a function definition statement, a jump statement or an empty statement.

Semantics: The expression expr1 is evaluated. While the boolean expression cond_expr is evaluated to true, the statement is executed and the expression expr2 is evaluated.
Note that a C-for statement does not return any atom
The following table presents several examples of C-for statements:


\section*{collection-for statement}

\section*{Syntax: for ( var in expr ) statement}
where cond_expr is a boolean expression, var denotes the name of a variable, expr is an expression of type collection and statement any statement: an expression statement, an iteration statement, a compound statement, a selection statement, a function definition statement, a jump statement or an empty statement.

Semantics: For each element in the collection denoted by expr, the variable var is assigned to this element and the statement is executed. Note that a collection-for statement does not return any atom

The following table presents several examples of collection-for statements:
```

for (x in list(1, 2, 3)) a += x;
increments a with 1, 2 and 3
for (x in (select Person)) names += x.name;
concatenates all the person names
for (x in (select Person.name = "john")) if (x.age < 10 || x.spouse.age
< 10) throw "cannot mary children!!";
a moralistic example
for (x in 1) doit();
raises an error: boolean expected, got integer

```

\subsection*{5.33 Jump Statements}

There are two jump statements based on the keywords break and return.

\section*{break Statement}

The syntax for break statement is:
break statement
where statement is an optional integer expression statement indicating the break level. If not specified, the break level is 1. The break statement may appear only within an iteration statement: while, do/while, for statements.
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|c|}{break Statement Examples} \\
\hline Statement & Comments \\
\hline break; & \begin{tabular}{l}
raises an error: break operator \\
<<break; >> : level 1 is too deep.
\end{tabular} \\
\hline break 3; & \begin{tabular}{ll} 
raises an error: & break operator \\
<<break; >> : level 3 is too deep.
\end{tabular} \\
\hline for ( x : \(=0\); ; \(\mathrm{x}++\) ) if ( \(\mathrm{x}==30\) ) break; & break the loop when x is equal to 30 \\
\hline ```
while (true) {x++; if (!check(x))
break;}
``` & break the loop when check(x) is not true \\
\hline ```
while (true) {x++; if (!check(x))
break 2;}
``` & \begin{tabular}{l}
raises an error: break operator \\
<<break; >> : level 2 is too deep.
\end{tabular} \\
\hline ```
while (true) {x++; while (x < 100)
if (!check(&x)) break 2;}
``` & break the two loop when check(\&x) is not true \\
\hline
\end{tabular}

\section*{return Statement}

The syntax for return statement is: return [statement]
where statement is an optional expression statement indicating the atom to return. The atom to return may be of any type. If not specified, no atom is returned. The return statement may appear only within a function definition statement.
\begin{tabular}{|c|c|}
\hline \multicolumn{2}{|c|}{return Statement Examples} \\
\hline Statement & Comments \\
\hline return; & raises an error: return operator <<return; >> : return must be performed in a function \\
\hline ```
function fact(n) {return n *
fact(n-1);}
``` & ok. \\
\hline ```
function f(x) {if (g(x)) return x+1;
return x+2;}
``` & ok. \\
\hline ```
function f(b) {if (b) return list(1,
2, 3); return bag(1);}
``` & ok. \\
\hline ```
function f(b) {if (b) return list(1,
2, 3); return bag(1);}
``` & ok. \\
\hline define as \(\mathrm{f}(\mathrm{x})\) (if (x) return x ) & raises an error: syntax error \\
\hline
\end{tabular}

\subsection*{5.34 Function Definition Statements}

As introduced previously, OQL supports functions. There are two sorts of functions definition syntax: function definition expression and function definition statements. The first ones, exposed in Section 5.26 can contain only one expressions. That means that they cannot include neither selection, neither statements, neither jump statements. Furthermore, as only one expression is allowed, functions that need several expressions, one must use the comma sequencing operaot to seperate the expressions, thus making this code not always readable.
We introduce here the more general form of function definitions which overrides the limitations of the previous form. The general form of function defintion statements is: function identifier ([arglist]) compound_statement
1. identifier denotes any valid OQL identifier, except a keyword
2. arglist is an optional comma-separated list of identifiers optionally followed, for default arguments, by a "?" and an expr, for instance:
(var1, var2, var3? expr, var4? expr)
3. compound_statement is a optionnal semicolon-separated list of statements surrounded by braces.

For instance:
```

function f(x, y, z ? oql\$maxint) {
if (x > y)
throw "error \#1";
return x - y * 2 / z;
}

```

\section*{Argument Types/Return Type}

Functions are not typed. That means that neither the return type nor the argument types may be given. It is why there is no function overloading mechanisms. To take benefit of the overloading mechanisms, one must use methods.
Nevertheless, it is possible to add type checking by using library functions such as is_int, is_string... combined with the assert or the assert_msg library functions. For instance, to check that the first argument is an integer and the second one a collection:
```

function doit(n, coll) {
assert_msg(is_int(n), "doit: argument \#1: integer expected");
assert_msg(is_coll(coll), "doit: argument \#2: collection expected");
// body of the function
}

```

The assert_msg check that its first argument is equal to true, otherwiser an exception containing the second argument string is thrown:
doit(1, list(1, 2, 3)); // ok
doit(1.2, list(1)); // raises the error:
// assertion failed: 'doit: argument \#1: integer expected'

\section*{Arguments in, out and inout}

Furthermore, one cannot specify that an argument is an input argument (in), an output argument (out) or an input/output argument (inout). In a function call, expressions and variables are always passed by value not by reference, this means that the call to "perform \((\mathrm{x}, \mathrm{y})\) " cannot modify neither x nor y . (In fact, yes it can! It is explained below. But forget it for now).
So, to modify variable through a function call, one needs to give the reference (or address) of this variable, not its value. In this case, the function must execute specific code dealing with address variables instead of their values.
The refof operator, introduced in a previous section, gives the reference of an identifier. Remember that the expression refof x returns the identifier x . To make a function call using references one must do: swap(refof x , refof y ) or the equivalent more compact form \(\operatorname{swap}(\& x, \& y\) ).
Contrary to C++, reference manipulation is not transparent in OQL: to access the value of a reference, one must use the valof operator (i.e. * operator). The swap function which swaps its two inout arguments has already been introduced:
```

function swap(rx, ry) {
v := *rx;
*rx := *ry;
*ry := v;
}

```

The arguments have been prefixed by \(r\) to indicate that they are references. So, the function call swap (\&x, \&y) will swap properly the variables x and y .
One can add type checking in the swap function, as follows:
```

function swap(rx, ry) {
assert_msg(is_ident(rx), "swap: argument \#1 identifier expected");
assert_msg(is_ident(ry), "swap: argument \#2 identifier expected");
v := *rx;
*rx := *ry;
*ry := v;
}

```

\section*{Return Value}

By default, a statement-oriented function returns no atom. To make a function returning an atom, one must use the return statement previously introduced. As a function has no specified returned type, it may contained several return statements returning atom of different types:
```

function perform(x) {
if (x == 1)
return "hello";
if (x == 2)
return list(1, 2, 3) + list(4, 20);
if (x == 3)
return 2;
if (x == 4)
return 'a';
}
alpha := perform(1); // alpha is equal to "hello"
alpha := perform(3); // alpha is equal to 2
alpha := perform(8); // alpha is equal to nil

```

\section*{Default Arguments}

OQL provides support for default arguments in a function definition statement. The syntax for a default argument is: "var ? expr" or "var := expr".
As in C and \(\mathrm{C}++\), the arguments with a default value must not followed by any argument with default values. For instance, function \(f(x, y, z:=\) "alpha") is valid while function \(f(x, y, z:=\) "alpha", t) is not valid.

\section*{Unval Arguments}

Sometimes, it is interesting to prevent the evaluation of some input arguments. For instance, let the function if_then_else which takes three arguments:
1. cond: a boolean expression,
2. then_expr: expression of any type; is evaluated and returned if and only if the condition is evaluated to true
3. else_expr: expression of any type; is evaluated and returned if and only if the condition is evaluated to false

It is clear that the following function definition:
```

function if_then_else(cond, then_expr, else_expr) {
if (cond)
return then_expr;
return else_expr;

```
is not correct as, although it returns the correct expression, the then_expr and the else_expr will be evaluated. For instance, if_then_else \((x<10,:: a:=2,: b:=3)\) will return 2 if \(x\) is less than 10 , otherwise it will return 3 , but in any case, a will be assigned to 2 and b will be assigned to 3 .
So, one needs a way to tell the interpreter that we do not want to evaluate the second and the third argument. The special character | before an argument means that this argument must not be evaluated. In this case, this argument is substitued by the string representation of the expression. For instance, let the function if_then_else:
```

function if_then_else(cond, |then_expr, |else_expr) {
// ...
}

```
when performing the call "if_then_else(x < 10, ::a := 2, ::b := 3)":
1. the value of cond in the body of the function will be true or false,
2. the value of then_expr in the body of the function will be "::a:=2"
3. the value of else_expr in the body of the function will be "::b:=3"

The correct implementation of this function is as follows:
```

function if_then_else(cond, |then_expr, lelse_expr) {
if (cond)
return eval then_expr;
return eval else_expr;
}

```

\section*{Scope of Variables}

In the body of a function defintion, every variable on the left side of an assignment has a local scope except if this variable is prefixed by the global scope operator \(::\). That means, that after the following statement sequence:
```

a := 2;
function doit() {
a := 1;
}

```
the variable a is still equal to 2 . While after:

\section*{Recursivity}
```

a := 2;

```
function doit() \{
    ::a := 1;
\}
the variable a is equal to 1.

\section*{Particularity}

One can define a statement-oriented function inside the body of a statement-oriented function, for instance:
```

function compute(amount_in_euro, usdollar_per_euro) {
function euro2usdollar(euro, usd ? usdollar_per_euro) {
return euro * usd;
}
x := euro2usdollar(euro * 1.24);
x += euro2usdollar(1000);
return x * .120;
}

```

Note that the function defined in the body of the function compute has a global scope, that means that after one execution of compute the function is available at the global level of the OQL session. It is possible, that in a future version, the functions defined in the body of a function definition will have a local scope.

\section*{The oql\$functions Variable}

The oql\$functions value is a list whose elements are the name of all the OQL functions defined in the current OQL session. Each you add a user function, this variable is updated. At the beginning of a session, the value of textttoql\$functions is:
list(is_int, is_char, is_double, is_string, is_oid, is_num, is_bool, is_bag, is_set, is_array, is_list, is_coll, is_struct, is_empty, void, assert, assert_msg, min, max, first, last, cdr, count, interval, sum, avg, is_in, distinct, flatten, flatten1, tolower, toupper, tocap, toset, tolist, tobag, toarray, listtoset, bagtoset, arraytoset, listtobag, settobag, arraytobag, bagtolist, settolist, arraytolist, bagtoarray, settoarray, listtoarray, strlen, substring, forone, forall, delete_from, get_from, ifempty, null_ifempty, getone)

For instance, to put the complete definition of all these functions into the variable functionString:
```

functionString := "";
for (x in oql\$functions)
functionString += "FUNCTION " x + ": " + bodyof x + "\n";

```

The next section provides a few statement-oriented and expression-oriented function definitions.

\section*{6 Quick Reference Manual}

This OQL quick reference manual presents consise information about the builtin and library functions and methods, the special variables, and the eyedboql tool. The standard library source code is presented and it provides a quick reference card containing all the language constructs.

\subsection*{6.1 Builtin and Library Functions and Methods}

OQL provides a few builtin and library functions. The builtin functions are written in C++ and cannot be overriden while the library functions are written in OQL and may be redefined by the user. The code for the library functions can be found in the section The Standard Library Package. The EyEDB system classes object, database, connection and EyeDB contains builtin class and instance methods that can be accessed from OQL. Some of these methods are briefly introduced in this section.

\section*{Type Predicat Functions}
is_int ( x ) : returns true if x is an int
is_double(x) : returns true if \(x\) is a double; otherwise, returns false
is_string ( \(x\) ) : returns true if \(x\) is a string; otherwise, returns false
is_oid ( \(x\) ) : returns true if \(x\) is an oid; otherwise, returns false
is_num ( \(x\) ) : returns true if \(x\) is an number (int, float or char); otherwise, returns false
is_bool ( \(x\) ) : returns true if \(x\) is a bool; otherwise, returns false
is_bag ( \(x\) ) : returns true if \(x\) is a bag; otherwise, returns false
is_set (x) : returns true if \(x\) is a set; otherwise, returns false
is_array ( \(x\) ) : returns true if \(x\) is a array; otherwise, returns false
is_list ( x ) : returns true if x is a list; otherwise, returns false
is_coll( \(x\) ) : returns true if \(x\) is a collection; otherwise, returns false
is_struct ( \(x\) ) : returns true if \(x\) is a struct; otherwise, returns false
is_empty \((x)\) : returns true if \(x\) is nil; otherwise, returns false

\section*{Collection Conversion Functions}

The collection conversion functions take one collection argument and convert this collection to another collection type and returns the converted collection.
```

toset(coll) : converts coll to a set
tolist(coll) : converts coll to a list
tobag(coll) : converts coll to a bag
toarray(coll) : converts coll to a array
listtoset(coll) : checks that coll is a list then converts coll to a set
bagtoset(coll) : checks that coll is a bag then converts coll to a set
arraytoset(coll) : checks that coll is a array then converts coll to a set
listtobag(coll) : checks that coll is a list then converts coll to a bag
settobag(coll) : checks that coll is a set then converts coll to a bag
arraytobag(coll) : checks that coll is a array then converts coll to a bag
bagtolist(coll) : checks that coll is a bag then converts coll to a list
settolist(coll) : checks that coll is a set then converts coll to a list
arraytolist(coll) : checks that coll is a array then converts coll to a list
bagtoarray(coll) : checks that coll is a bag then converts coll to a array
settoarray(coll) : checks that coll is a set then converts coll to a array
listtoarray(coll) : checks that coll is a set then converts coll to a array

```

\section*{Sort Functions}

These functions are used to sort collection of sortable atom of homogeneous types: int, char, float or string.
```

sort(coll) : coll must a be a collection of homogeneous sortable atoms;
sorts and returns this collection
rsort(coll) : coll must a be a collection of homogeneous sortable atoms;
reverse sorts and returns this collection
coll must a be a collection of list or array
of homogeneous sortable atoms;
idx must be of int type;

```
sorts the collection of collections according to the \#idx element of the inner collection
risort (coll, idx) : same as previous function, but perform a reverse sort

\section*{Collection Miscelleanous Functions}
\begin{tabular}{|c|c|}
\hline first(coll) & : returns the first element of coll \\
\hline car (coll) & : returns the first element of coll \\
\hline last(coll) & : returns the last element of coll \\
\hline cdr (coll) & : returns all elements of coll but the first \\
\hline \(\operatorname{getn}(\operatorname{coll}, \mathrm{n})\) & : returns at most n elements of coll \\
\hline count (coll) & : returns the count of elements of coll identical to coll[!], but less efficient \\
\hline sum (coll) & : returns the sum of the numbers of coll \\
\hline avg (coll) & : returns the float average of the numbers coll \\
\hline distinct(coll) & : eliminates duplicates of coll \\
\hline flatten(coll) & : recursive flattening of coll \\
\hline flatten1 (coll) & : one level flattening of coll \\
\hline min (coll) & : returns the minimal number of coll \\
\hline \(\max\) (coll) & : returns the maximal number of coll \\
\hline forone(coll, f, data) & if \(f(e\), data) for one element e of coll, returns true; otherwise returns false; \\
\hline forall (coll, f, data) & if \(f(e\), data) for all element e of coll, returns true; otherwiser returns false; \\
\hline
\end{tabular}

\section*{String Function Functions}
\begin{tabular}{ll} 
tolower(str) & \(:\) converts (and returns) string str into lowercase \\
toupper(str) & \(:\) converts (and returns) string str into uppercase \\
tocap(str) & \(:\) converts the first character and each character following \\
& a - of str into an uppercase \\
strlen(str) & \(:\) returns the length of str; \\
& same as str [!], but less efficient \\
substring(str, from, len) : returns the sub-string of str \\
& from the \#from to the \\
& from+len characters; \\
& same as str[from:from+len] but less efficient
\end{tabular}

\section*{Query Functions}
delete_from(class) : deletes all the instances of a given class
get_from(class) : gets all the instances of a given class

\section*{Useful Functions}
\begin{tabular}{ll} 
assert (cond) & : throws an exception is cond is not true \\
assert msg(cond, msg) & : throws the exception message msg if cond is not true \\
interval (from, to) & : returns a list composed of the number from from to to
\end{tabular}

\section*{Native Methods of the Class object}

The native methods of the class object allows us to perform a few action such as getting the oid of an instance get0id(), getting the database of an instance getDatabase() or converts the instance to its string representation toString(). For instance, to apply this last method to the first Person instance: first(select Person).toString().

All the native methods of the class object are instance methods.
```

oid getOid() : returns the oid of the object
string toString() : returns the string representation of the object
database *getDatabase() : returns the database instance of the object
void setDatabase(in database *) : changes the database of the object

```
```

void store() : stores the object in the database
object *clone() : clones the object; returns the clone
int getCTime() : returns the creation time of the object (seconds from 1/1/1970)
int getMTime() : returns the last modification time of the object
string getStringCTime() : returns the string representation of the creation time of the object
string getStringMTime() : returns the string representation of the creation time of the object
bool isRemoved()
bool isModify() : returns true if the object is modified; false otherwise

```

\section*{Native Methods of the Class connection}

All the native methods of the class connection are instance methods. an object obtained using the new operator, They can be applied on a database object that can be either the current connection oql\$db->getConnection() either an object obtained using the new operator, for instance: new <> connection().
void open() : opens a new connection with default host and port
void open(in string host, piensstning cpontètion using host and port
void close() : closes the connection

\section*{Native Methods of the Class database}

The following methods are the instance methods of the class database: They can be applied on a database object that can be either oql\$db either an object obtained using the new operator, for instance: new <> database(name : "foo").
```

void open(in connection *conn,
in int mode) : opens a new database using the connection
conn and the open flag mode mode
void open(in connection *conn,
in int mode,
in string userauth,
in string passwdauth)
void close()
connection *getConnection()
int getOpenMode()
int getVersionNumber()
string getVersion() : returns the string version of the database
void removeObject(in oid) : removes the object whose oid is given
void transactionBegin() : begins a new transaction
void transactionBegin(in string mode) : begins a new transaction in mode mode
void transactionCommit() : commits the current transaction
void transactionAbort() : abort the current transaction
bool isInTransaction() : returns true if a transaction is in progress;
false otherwiser

```

\section*{Native Methods of the Class EyeDB}

All the native methods of the class EyeDB are class methods.
string getConfigValue(in string s) : returns the string value of the configuration variable s
int getVersionNumber () : returns the EYEDB current version number
string getVersion() : returns the EYEDB current stringversion
string getArchitecture() : returns the architecture of the current server
string getDefaultCompiler () : returns the C++ compiler used to compile the current server

\subsection*{6.2 Special Variables}

The following variables are predefined or have a special meaning:
oql\$variables : list containing the name of all variables
oql\$functions : list containing the name of all functions
oql\$result : the result atom of the last statement
oql\$db : object atom instance of the class database denoting the current database
oql\$minint : the minimal integer 0x8000000000000000
oql\$maxint : the maximal integer 0x7FFFFFFFFFFFFFFF
oql\$minfloat : the minimal float 4.94065645841246544e-324
oql\$maxfloat : the maximal float \(1.79769313486231570 \mathrm{e}+308\)

\subsection*{6.3 The eyedboql Tool}
eyedboql is a tool that allows you to execute interactively OQL statements. This tool is similar to the Oracle sqlplus and Sybase isql well known tools.

\section*{Running eyedboql}

The simplest way to run eyedboql is to type eyedboql as follows (assuming that \(\$\) is your shell prompt):
```

\$ eyedboql
Welcome to eyedboql.
Type '\help' to display the command list.
Type '\copyright' to display the copyright.
?

```

The string "? " is the default prompt for eyedboql.
In an eyedboql session, anything you type is interpreted as an OQL construct (or a part of an OQL construct), Nevertheless, if the first word of an input line begins with the escape character "
", this word is interpreted as an eyedboql internal command. There are about 30 internal commands, but you need to know only of few of them to use eyedboql.

The purpose of the main internal commands is:
- to create or delete databases,
- to open a database,
- to begin, commit or abort a transaction,
- to set the current user and password,
- to execute the contents of a file,
- to display the string representation of an object,
- to display the HTML representation of an object in a WEB browser,
- to change the prompts and the escape character.

\section*{Executing Statements}
eyedboql allows us to execute OQL statements in an interactive way. For instance, to perform the addition \(1+3\) :
```

\$ eyedboql
Welcome to eyedboql.
Type '\help' to display the command list.
Type '\copyright' to display the copyright.
? 1+3;
= 4
?

```

The string "= " preceedes the result atom (if any) of your statement; in the current example, the result atom is the evaluation of the expression statement \(1+3\);

\section*{Execution Process}

Each complete statement typed is sent to the OQL interpreter, A complete statement has a special meaning: it is any sequence of characters:
\(a\). that end with a semi-colon and
b. which parenthesis are balanced and
c. which brakets are balanced and
\(d\). which braces are balanced.
While the statement is not complete, eyedboql prompts the "second prompt" (" \(\gg\) " by default) after each newline.
Once the statement is complete, it is sent to the OQL interpreter, then the atom result is display (if any) after the string "= " and the main prompt "? " is displayed again.

For instance, the input sequence of characters " \(1+\) newline 3 newline; " gives:
```

? 1+
>> 3
>> ;
= 4

```
while the input sequence " \(\{\mathrm{a}:=1+3\); newlinec \(:=2+94\); newline \(\mathrm{d}:=\mathrm{a}+\mathrm{c}\}\) " gives:
```

? { a := 1+3;
>> c := 2+94;
>> d := a+c}

```
?

Note that no "= result atom" is echoed because a compound statement does not return any atom.
This last example:
```

? while (true) {
>> a++;
>> b++;
>> }
>>
>> ;
?

```
shows the necessity of typing a semicolon after the while statement although a while statement does not need to end by a semi-colom in the OQL specifications.

\section*{Getting Started}

By default in an eyedboql session, the database EYEDBDBM is opened in read-only mode. To use another database, one must use either the command line option -db either the
open internal command.
To start to play with eyedboql one needs to know the following five internal commands:
1.
open database [rw|ro] local trsless: opens the database in read-write (rw) or read-only (ro or no option) mode. If local is there, database is opened in a local mode. If trsless is there database is opened in transaction-less mode. For instance "
open foo rw" opens the database foo is read-write mode.
2.
print [sequence of oids] other options: if a sequence of oids is given: loads the object corresponding to each oid and displays its string representation,
if no sequence of oids is given: loads the object corresponding to each oid returned by the last statement and displays its string representation.
For instance, "select Person;" will return a bag containing the oid of each Person instance.
The internal command "
print" typed after that will loads and displays of the corresponding Person instances.
The other options are one or more space-separated of the followings:
\begin{tabular}{ll} 
full & \(:\) loads and displays object using the full recursive mode \\
ctime & : displays the creation time of the object \\
mtime & \(:\) displays the last modification time of the object \\
contents & \(:\) displays the contents of collections \\
native & \(:\) displays the native attributes \\
all & \(:\) means "ctime mtime contents native"
\end{tabular}

For instance "
print full contents" will load and display the objects and their collection contents in a full recursive mode.
3.
commit: commits the current transaction
4.
abort: aborts the current transaction
5.
quit or \({ }^{\wedge} D\) : quits the current eyedboql session
Note that a transaction is started automatically before the first complete statement of the sesssion of before the complete statement immediately following
commit or
abort internal command.
Here is a commented example showing the use of these internal commands:
```

run eyedboql:
\$ eyedboql
Welcome to eyedboql.
Type '\help' to display the command list.
Type '\copyright' to display the copyright.

```
open the database person in read-write mode:
? \open person rw
get the first person whose name is "john" and display it:
? john := first(select Person.name = "john");
= 66373.12.4008447:oid
? \print
66373.12.4008447:oid Person = \{
name = "john";
age \(=32\);
addr Address = \{
            street = "clichy";
            town = "Paris";
            country = NULL;
\};
cstate \(=\) Sir;
*spouse 66891.12.2738687:oid;
cars set<Car*> = set \{
name \(="\) ";
count \(=4\);
dimension \(=1\);
reference \(=\) true;
magorder \(=4\);
\};
children array<Person*> = array \{
name \(=" " ;\)
count \(=0\);
range \(\quad=[0,0[\);
dimension \(=1\);
reference \(=\) true;
magorder \(=4\);
\};
\(\mathrm{x}=\) NULL;
\};
change the name of john to "JOHNNY":
? john.name := "JOHNNY";
= "JOHNNY"
retrieve the person whose name is "JOHNNY" and compares it to john using assert : all is fine, no error is raised!
? assert(john = first(select Person.name = "JOHNNY"));
abort the transaction and look for the person whose name is "JOHNNY": no person is returned! this is ok as the transaction was aborted:
? \abort
? select Person.name = "JOHNNY";
= list()
change the name of john to "JOHNNY" again and commit the transaction:
? john.name := "JOHNNY";
= "JOHNNY"
? \commit
then retrieve again the person whose name is "JOHNNY" and compare it to john using assert: all is fine, no error is raised!
? assert(john = first(select Person.name = "JOHNNY"));
quit eyedboql session
? \quit
\$
We are going to conclude this section by this important note:
as introduced previously, the current transaction will be committed (resp. aborted) by a
commit (resp.
abort) command.
If you quit eyedboql before committing (resp. aborting) the transaction, it will be automatically aborted, so all your changes will be lost.
This is the default behaviour. This behaviour can be changed by using the commitonclose internal command.

\section*{Command Line Options}

The eyedboql command line options are as follows:

\section*{Program Options:}
-d <name>, -database=<name \({ }_{i}\)
-r, -read
-w, -read-write
-s, -strict-read
-l, -local
-c <command>, -command=<command>
-p, -print
-full
-commit
-i, -interact
-e, - echo
-admin
-h, -help
<file \({ }_{i}\)

\section*{Common Options:}
-U <user>|@, -user=<user>|@
-P [<passwd>], -passwd \([=<\) passwd \(>]\)
-host \(=<\) host>
-port \(=<\) port \(>\)
-inet
-dbm=<dbmfile>
-conf=<conffile>
-logdev=<logfile>
-logmask=<mask>
-logdate=on|off
-logdate \(=\) on \(\mid\) off

Database name
Open database in read mode
Open database in read/write mode
Open database in strict read mode
Open database in local mode
OQL command to execute
Display all the objects loaded
Full recursive mode is used to display objects
Commits the current transaction on close
Enter interpreter after executing file or commands
Echo each command
Open database in admin mode
Display this message
File(s) to execute

User name
Password
eyedbd host
eyedbd port
Use the tcp_port variable if port is not set
EYEDBDBM database file
Configuration file
Output log file
Output log mask
Control date display in output log
\begin{tabular}{ll}
-logtimer=on loff & Control timer display in output log \\
-logpid=on |off & Control pid display in output \(\log\) \\
-logprog=on |off & Control progname display in output log \\
-error-policy \(=<\) value> & Control error policy: status lexception |abort | stop l echo \\
-trans-def-mag=<magorder> & Default transaction magnitude order \\
-arch & Display the client architecture \\
-v, -version & Display the version \\
-help-eyedb-options & Display this message
\end{tabular}

For instance, to execute the statement "delete_from(Person)" on the database person:
```

\$ eyedboql -d person -w -c "delete_from(Person)"

```
\$
To execute the command "persons \(:=\) (select Person)" and then enter the interactive mode of eyedboql:
```

\$ eyedboql -d person -w -c "persons := (select Person)" -i
Welcome to eyedboql.
Type '\help' to display the command list.
Type '\copyright' to display the copyright.
?

```

To execute the file mylib.oql:
\$ eyedboql -d person -w mylib.oql

\subsection*{6.4 The Standard Library Package}

The stdlib.oql file contains a few basic library functions. It can be found in the directory libdir/eyedb/oql. It is automatically imported when opening an OQL session. The following code, extracted from this file, provides an interesting way to understand OQL.
```

//
// minimal and maximal integer values
//
oql$maxint := 0x7FFFFFFFFFFFFFFF;
oql$minint := 0x8000000000000000;
nulloid := 0:0:0:oid;
NULLOID := 0:0:0:oid;
//
// type predicat functions
//
define is_int(x) as (typeof x == "integer");
define is_char(x) as (typeof x == "char");
define is_float(x) as (typeof x == "float");
define is_string(x) as (typeof x == "string");
define is_oid(x) as (typeof x == "oid");
define is_object(x) as (typeof x == "object");
define is_num(x) as (is_int(x) || is_float(x) || is_char(x));
define is_bool(x) as (typeof x == "bool");
define is_bag(x) as (typeof x == "bag");
define is_set(x) as (typeof x == "set");
define is_array(x) as (typeof x == "array");
define is_list(x) as (typeof x == "list");
define is_coll(x) as (is_list(x) || is_array(x) || is_set(x) || is_bag(x));
define is_struct(x) as (typeof x == "struct");
define is_empty(x) as (typeof x == "nil");
//
// void(x): evaluates argument and returns nil
//

```
```

define void(x) as (x, nil);
function assert(|cond) {
r := eval cond;
if (!r)
throw "assertion failed: '" + cond + "'";
}
function assert_msg(|cond, msg) {
r := eval cond;
if (!r)
throw "assertion failed: '" + msg + ")";
}
//
// min(l): returns the minimal integer in a collection
//
function min(l) {
msg := "function min(" + (string l) + "): ";
if (!is_coll(l))
throw msg + "collection expected";
m := oql$maxint;
        for (x in l) {
        if (x != null) {
            if (!is_num(x))
                throw (msg + "numeric expected");
            if (x < m)
                m := x;
        }
    }
    return m;
}
//
// max(l): returns the maximal integer in a collection
//
function max(l) {
    msg := "function max(" + (string l) + "): ";
    if (!is_coll(l))
        throw msg + "collection expected";
    m := oql$minint;
for (x in l) {
if (x != null) {
if (!is_num(x))
throw (msg + "numeric expected");
if (x > m)
m := x;
}
}
return m;
}
//
// first(l): returns the first element in a list or array
//

```
    if (!is_coll(l)) // if (!is_list(l) && !is_array(l))
        throw "function first: collection expected";
    start := 0;
    f := nil;
    for (x in l)
        if (start == 0) {
            start := 1;
            f := x;
            break;
        }
    return f;
}
car := &first;
//
// last(l): returns the last element in a list or array
//
function last(l) {
    if (!is_coll(l)) // if (!is_list(l) && !is_array(l))
        throw "function last: list or array expected";
    f := nil;
    for (x in 1)
        f := x;
    return f;
}
//
// cdr(l): returns all elements in a collection except the first one
//
function cdr(l) {
    if (!is_coll(l)) // if (!is_list(l) && !is_array(l))
        throw "function cdr: list or array expected";
    r := list();
    n := 0;
    for (x in l) {
        if (n != 0)
            r += x;
        n++;
    }
    return r;
}
//
// getn(l, n): returns all elements in a collection
//
function getn(l, n) {
    rl := list();
    cnt := 0;
    for (x in l) {
        if (cnt++ >= n)
            break;
        rl += x;
    }
    return rl;
```

```
}
//
// getrange(l, f, nb): returns all elements in a collection from element from
//
// identical to l[f:f+nb]
//
function getrange(l, f, nb) {
    if (!is_list(l) && !is_array(l))
        throw "function getrange: list or array expected";
        rl := list();
        cnt := 0;
        max := f + nb;
        for (x in l) {
            if (cnt >= max)
                break;
            if (cnt >= f)
                rl += x;
            cnt++;
    }
    return rl;
}
//
// count(l): returns element count of a collection
//
function count(l) {
        if (typeof l == "nil")
            return 0;
        if (!is_coll(1))
            throw "function count: collection expected, got " + typeof(1);
        return 1[!];
}
//
// interval(x, y): constructs an integer list bounded by 'x' and 'y'
//
function interval(x, y) {
    n := x-1;
    l := list();
    while (n++ < y)
        l += n;
    return l;
}
//
// sum(l): returns the sum of collection elements
//
function sum(1) {
    if (!is_coll(1))
            throw "function sum: collection expected";
    n := 0;
    for (x in l)
```

```
        n += x;
    return n;
}
//
// avg(l): returns the average of collection elements
//
function avg(1) {
    if (!is_coll(1))
        throw "function avg: collection expected";
    return float(sum(1))/count(l);
}
//
// is_in(l, z): returns true in element 'z' is in collection 'l'
//
function is_in(l, z) {
    for (x in l)
        if (x == z)
            return true;
    return false;
}
//
// distinct(l): returns distinct elements in a collection
//
function distinct(l) {
    if (is_list(1)) ll := list();
    else if (is_bag(l)) ll := bag();
    else if (is_array(l)) ll := array();
    else if (is_set(l)) ll := set();
    else throw "function distinct: collection expected";
    for (x in l)
        if (!is_in(ll, x))
            ll += x;
    return ll;
}
//
// flatten(l): full recursive flatten function
//
function flatten(l) {
    if (!is_coll(l))
        return l;
    ll := list();
    for (x in l)
        if (is_coll(x))
            ll += flatten(x);
        else
            ll += x;
        return 1l;
}
//
// flatten1(1): 1-level recursive flatten function
```

```
//
```

function flatten1(1) \{
if (!is_coll(l))
return l;
ll := list();
for ( x in 1 )
ll += x ;
return ll;
\}
//
// tolower(s): returns lower case string
//
function tolower(s) \{
$\mathrm{n}:=0$;
x := "";
delta := 'a' - 'A';
while (s[n] != '\000') \{
if ( $\mathrm{s}[\mathrm{n}]>=$ 'A' \&\& $\mathrm{s}[\mathrm{n}]<=$ ' Z ')
$\mathrm{x}+=\operatorname{string}(\operatorname{char}(\mathrm{s}[\mathrm{n}]+\operatorname{delta}))$;
else
$\mathrm{x}+=\operatorname{string}(\mathrm{s}[\mathrm{n}])$;
n++;
\}
return x ;
\}
//
// toupper(s): returns upper cased string
//
function toupper (s) \{
$\mathrm{n}:=0$;
x := "";
delta := 'A' - 'a';
while (s[n] != '\000') \{
if ( $\mathrm{s}[\mathrm{n}]>=$ 'a' \&\& $\mathrm{s}[\mathrm{n}]<=$ ' z ')
$\mathrm{x}+=\operatorname{string}(\operatorname{char}(\mathrm{s}[\mathrm{n}]+\operatorname{delta}))$;
else $x+=\operatorname{string}(s[n])$;
n++;
\}
return x ;
\}
//
// tocap(s): returns capitalized word string
//
function tocap(s) \{

```
n := 1;
x := "";
delta := 'A' - 'a';
```

```
    s := tolower(s);
    if (s[0] >= 'a' && s[0] <= 'z')
        x += string(char(s[0] + delta));
    while (s[n] != '\000') {
        if (s[n] == ' _')
            x += string(char(s[++n] + delta));
        else
            x += string(s[n]);
        n++;
    }
    return x;
}
//
// Collection Conversion Functions
//
//
// General Conversion Functions
//
function toset(l) {
    if (!is_coll(l))
        throw ("function toset: collection expected, got " + typeof(l));
    if (!is_set(l)) {
        s := set();
        for (x in l)
            s += x;
        return s;
    }
    return 1;
}
function tolist(1) {
    if (!is_coll(l))
        throw ("function tolist: collection expected, got " + typeof(l));
    if (!is_list(l)) {
        s := list();
        for (x in l)
            s += x;
        return s;
    }
    return 1;
}
function tobag(l) {
    if (!is_coll(1))
        throw ("function tobag: collection expected, got " + typeof(1));
    if (!is_bag(l)) {
        s := bag();
        for (x in l)
            s += x;
        return s;
    }
```


## return 1;

\}
function toarray(l) \{
if (!is_coll(l))
throw ("function toarray: collection expected, got " + typeof(1));
if (!is_array(l)) \{ s := array () ; for ( x in 1 )
s += $x$;
return s;
\}
return 1;
\}
//
// toset family Conversion Functions
//
function listtoset(l) \{
if (!is_list(l))
throw ("function listtoset: list expected, got " + typeof(1));
return toset(1);
\}
function bagtoset(l) \{
if (!is_bag(l))
throw ("function bagtoset: bag expected, got " + typeof(1));
return toset(1);
\}
function arraytoset(l) \{
if (!is_array(l))
throw ("function arraytoset: array expected, got " + typeof(1));
return toset(1);
\}
//
// tobag family Conversion Functions
//

```
function listtobag(l) {
    if (!is_list(l))
        throw ("function listtobag: list expected, got " + typeof(l));
    return tobag(l);
}
function settobag(l) {
    if (!is_set(l))
        throw ("function settobag: set expected, got " + typeof(1));
    return tobag(1);
}
function arraytobag(l) {
    if (!is_array(l))
        throw ("function arraytobag: array expected, got " + typeof(1));
    return tobag(l);
}
```

//

```
// tolist family Conversion Functions
//
function bagtolist(l) {
    if (!is_bag(1))
        throw ("function bagtolist: bag expected, got " + typeof(1));
    return tolist(1);
}
function settolist(l) {
    if (!is_set(l))
        throw ("function settolist: set expected, got " + typeof(1));
    return tolist(1);
}
function arraytolist(1) {
    if (!is_array(1))
        throw ("function arraytolist: array expected, got " + typeof(1));
    return tolist(1);
}
//
// toarray family Conversion Functions
//
function bagtoarray(l) {
    if (!is_bag(1))
        throw ("function bagtoarray: bag expected, got " + typeof(1));
    return toarray(l);
}
function settoarray(l) {
    if (!is_set(l))
            throw ("function settoarray: set expected, got " + typeof(1));
    return toarray(1);
}
function listtoarray(l) {
    if (!is_list(l))
        throw ("function listtoarray: list expected, got " + typeof(1));
    return toarray(1);
}
//
// strlen(s): same as s[!]
//
function strlen(s) {
    len := 0;
    while (s[len] != '\000')
        len++;
    return len;
}
//
// substring(str, f, len)
//
function substring(str, f, len) {
    s := "";
    n := 0;
    max := str[!] - f;
    while (n < len && n < max) {
```

```
        s += string(str[n+f]);
        n++;
    }
    return s;
}
//
// forone(l, fpred, data): returns true if and only if the function 'fpred'
// returns true for at least one element 'x' in list 'l'
//
function forone(l, fpred, data) {
    for (x in l)
        if (fpred(x, data)) return true;
    return false;
}
//
// forone(l, fpred, data): returns true if and only if the function 'fpred'
// returns true for all elements 'x' in list 'l'
//
function forall(l, fpred, data) {
    for (x in l)
        if (!fpred(x, data)) return false;
    return true;
}
//
// delete_from(cls): delete all instances of class 'cls'
//
function delete_from(|cls) {
    for (x in (eval "select " + cls))
        delete x;
}
//
// delete_(coll): delete contents of collection coll
//
function delete_(coll) {
    for (x in coll)
        delete x;
}
//
// get_from(cls): returns all instances of class 'cls'
//
function get_from(|cls) {
    eval "select " + cls;
}
//
// generates an unused global symbol
//
function gensym() {
    prefix := "::oql#_#_#";
    for (i := 0; ; i++) {
        varname := prefix + string(i);
```

```
        if (!(eval "isset " + varname)) {
            eval varname + " := 0";
            return ident(varname);
        }
    }
}
//
// expression-like for-each function
//
function foreach_expr(|x, |coll, lexpr, colltyp ? "list") {
    varname := "_#_#_R_#_#_";
    statement := "push " + varname + " := " + colltyp + "(); " +
        "for (" + x + " in " + coll + ") " +
        "{" + varname + " += " + expr + ";}" +
        "pop " + varname;
    return eval statement;
}
//
// expression-like for-C function
//
function for_expr(|start, |cond, lend, lexpr, colltyp ? "list") {
    varname := "_#_#_R_#_#_";
    statement := "push " + varname + " := " + colltyp + "(); " +
        "for (" + start + "; " + cond + "; " + end + ")" +
        "{" + varname + " += " + expr + ";}" +
        "pop " + varname;
    return eval statement;
}
//
// expression-like while-C function
//
function while_expr(|cond, lexpr, colltyp ? "list") {
    varname := "_#_#_R_#_#_";
    statement := "push " + varname + " := " + colltyp + "(); " +
        "while (" + cond + ")" +
        "{" + varname + " += " + expr + ";}" +
        "pop " + varname;
    return eval statement;
}
//
// expression-like do/while-C function
//
function do_while_expr(lexpr, |cond, colltyp ? "list") {
    varname := "_#_#_R_#_#_";
    statement := "push " + varname + " := " + colltyp + "(); " +
        "do {" + varname + " += " + expr + ";}" +
        "while (" + cond + ");" +
        "pop " + varname;
```

```
    return eval statement;
}
function extentof(|classname) {
    return (select one class.name = classname).extent;
}
function countof(|classname) {
    return (select one class.name = classname).extent.count;
}
function objectcount(db := oql$db) {
            objcnt := 0;
            db->transactionBegin();
            for (cl in (select <db> x from class x where
                    x.type != "system" and x.name !~ "<"))
            objcnt += cl.extent.count;
        db->transactionCommit();
        return objcnt;
}
function ifempty(x, y) {
    if (is_empty(x))
        return y;
    return x;
}
function null_ifempty(x) {
    return ifempty(x, null);
}
function getone(x) {
    if (is_empty(x))
        return null;
    return first(flatten(x));
}
//
// database and transaction management
//
function open_db(db_name_or_id, strmode, user := null, passwd := null) {
    if (strmode == "r")
    mode := DBREAD;
else if (strmode == "rw")
    mode := DBRW;
else if (strmode == "rlocal")
    mode := DBREAD|DBOPENLOCAL;
else if (strmode == "rwlocal")
    mode := DBRW|DBOPENLOCAL;
else
        throw "invalid open mode: r, rw, rlocal or rwlocal expected, got " +
            strmode;
if (is_int(db_name_or_id))
    db := new<> database(dbid : db_name_or_id);
else
    db := new<> database(dbname : db_name_or_id);
if (user == null)
    db.open(oql$db.getConnection(), mode);
```

```
    else
        db.open(oql$db.getConnection(), mode, user, passwd);
    return db;
}
function set_default(db) {
    db->setDefaultDatabase();
}
function begin(db := oql$db) {
    db->transactionBegin();
}
function begin_params(trsmode, lockmode, recovmode, magorder, ratioalrt, wait_timeout, db := oql$db) {
    db->transactionBegin(trsmode, lockmode, recovmode, magorder, ratioalrt, wait_timeout);
}
function commit(db := oql$db) {
    db->transactionCommit();
}
function abort(db := oql$db) {
    db->transactionAbort();
}
//
// miscellaneous
//
function print_function(f) {
    print "function " + (bodyof f) + "\n";
}
function print_functions() {
    cnt := 0;
    for (f in oql$functions) {
        if (cnt > 0) print "\n";
        print_function(f);
        cnt++;
    }
}
function print_variable(v) {
    print string(v) + " = " + string(eval string(v)) + ";\n";
}
function print_variables() {
    for (v in oql$variables) {
        print_variable(v);
        cnt++;
    }
}
function print_classes(system := false) {
    if (system)
        l := (select list(x, x.name) from class x order by x.name);
    else
        l := (select list(x, x.name) from class x where x.type = "user" and x.name !~ "<" order by x.name);
    for (c in l) {
        cls := c[0];
        clsname := c[1];
```

```
        print "class " + clsname;
        if (cls.parent != NULL && (system || cls.parent.type != "system"))
            print " extends " + cls.parent.name;
        print "\n";
    }
}
function print_obj(o, flags := 0) {
    print o->toString(flags);
}
function print_objs(l, flags := 0) {
    for (o in l)
        print_obj(o, flags);
}
//
// contents_ expression
//
function contents_(coll) {
    r := list();
    for (x in coll) {
        for (s in contents(x))
        r += s;
    }
    return r;
};
function println(s) {
    print(s+"\n");
}
function bench(|cmd) {
    t0 := time_stamp::local_time_stamp();
    r := eval cmd;
    t1 := time_stamp::local_time_stamp();
    us := t1->substract(t0).usecs;
    println("Elapsed time: " + string(us/1000.) + " ms");
    return r;
}
```


### 6.5 OQL Quick Reference Card

The following table presents all the OQL statements, expression types and the operators. For the operators common to $\mathrm{C}++$ and OQL, the precedence and associativity is the same.


| greater than <br> match regular expression <br> match regular expression case insensitive <br> not match regular expression <br> not match regular expression case insensitive match regular expression | $>=$ <br> !~ !~~ like | expr >= expr <br> expr ~ expr <br> expr ~~ expr <br> expr ! ~ expr <br> $\operatorname{expr}!{ }^{\sim}$ expr <br> expr like expr |
| :---: | :---: | :---: |
| Conditionnal Expressions |  |  |
| conditionnal expression | ? : | expr ? expr : expr |
| Expression Lists |  |  |
| comma sequencing | , | expr , expr |
| Array Expressions |  |  |
| subscripting interval subscriptiong | $\begin{aligned} & {[]} \\ & {[:]} \end{aligned}$ | $\begin{aligned} & \text { expr [expr ] } \\ & \text { expr [expr : expr ] } \\ & \hline \end{aligned}$ |
| Path Expressions |  |  |
| member selection member selection | -> | $\begin{aligned} & \text { expr . expr } \\ & \text { expr } \rightarrow \text { expr } \end{aligned}$ |
| Function Call |  |  |
| function call | () | expr (expr_list) |
| member selection | ```Method I ()``` | $\text { expr } \rightarrow \operatorname{expr} \text { (arglist) }$ |
| Eval/Unval Operators |  |  |
| eval <br> no eval | eval unval | eval expr unval expr |
| Identifier Expressions |  |  |
| scope <br> is set <br> unset <br> reference of <br> value of <br> value of <br> scope of <br> push onto symbol table <br> push onto symbol table and as- <br> sign <br> pop from symbol table | : : <br> isset <br> unset <br>  <br> refof <br> valof <br> scopeof <br> push <br> push <br> pop | :: identifier isset identifier unset identifier \& identifier refof identifier * identifier valof identifier scopeof identifier push identifier push expr <br> pop identifier |
| Set Expressions |  |  |
| union <br> intersection <br> except <br> include <br> include or equal <br> contain <br> contain or equal | union <br> intersect <br> except <br> < <br> <= <br> $>$ <br> $>=$ | ```expr union expr expr intersect expr expr except expr expr < expr expr <= expr expr > expr expr >= expr``` |
| Object Creation |  |  |
| new <br> new | new <br> new | [new] new_construct <br> new< opt_expr > new_construct |
| Object Deletion |  |  |
| delete | delete | delete expr |
| Collection Expressions |  |  |
| contents <br> is in <br> add to collection <br> suppress from collection <br> set element in or get element from an indexed collection <br> append to an indexed collection <br> empty collection <br> exists in collection <br> for all in collection | ```contents in add to suppress from [] append/to empty in for all``` | ```contents expr expr in expr add expr to expr suppress expr from expr expr [ expr ] append expr to expr empty expr exists identifier in expr : expr for all identifier in expr : expr``` |


| for some in collection | for | for < expr, expr > in expr : expr |
| :---: | :---: | :---: |
| Exception Expressions |  |  |
| throw exception | throw | throw expr |
| Function Definition |  |  |
| define function | define as | define identifier [arglist as expr |
| Conversion |  |  |
| string conversion | string | string ( expr ) |
| integer conversion | int | int ( expr ) |
| character conversion | cha | char ( expr) |
| float conversion | float | float ( expr ) |
| identifier conversion | ident | ident ( expr ) |
| oid conversion | oid | oid ( expr ) |
| Query Expressions |  |  |
| database query | select | ```select expr [from { expr [as] identifier} [where expr]] [order by {expr}] select expr [from {identifier in expr} [where expr]] [order by {expr}]``` |
| Type Information Expressions |  |  |
| class of typeof of | classof typeof | classof expr typeof expr |
| Miscellenaous Expressions |  |  |
| structure of body of | structof bodyof | structof expr bodyof expr |
| length pf import package | $\begin{aligned} & \text { [!] } \\ & \text { import } \end{aligned}$ | $\begin{aligned} & \text { expr [!] } \\ & \text { import expr } \end{aligned}$ |


[^0]:    - atomic expressions,
    - arithmetic expressions,
    - assignment expressions,
    - auto increment \& decrement expressions,
    - comparison expressions,
    - logical expressions,
    - conditional lists,
    - expression sequences,
    - array deferencing,
    - identifier expressions,
    - path expressions,
    - function call,
    - method invocation,
    - eval/unval operators,
    - set expressions,
    - object creation,
    - object deletion,
    - collection expressions,
    - exception expressions,
    - function definition expressions,
    - conversion expressions,
    - type information expressions
    - query expressions,
    - miscellenaous expressions

